

Computational Geometry

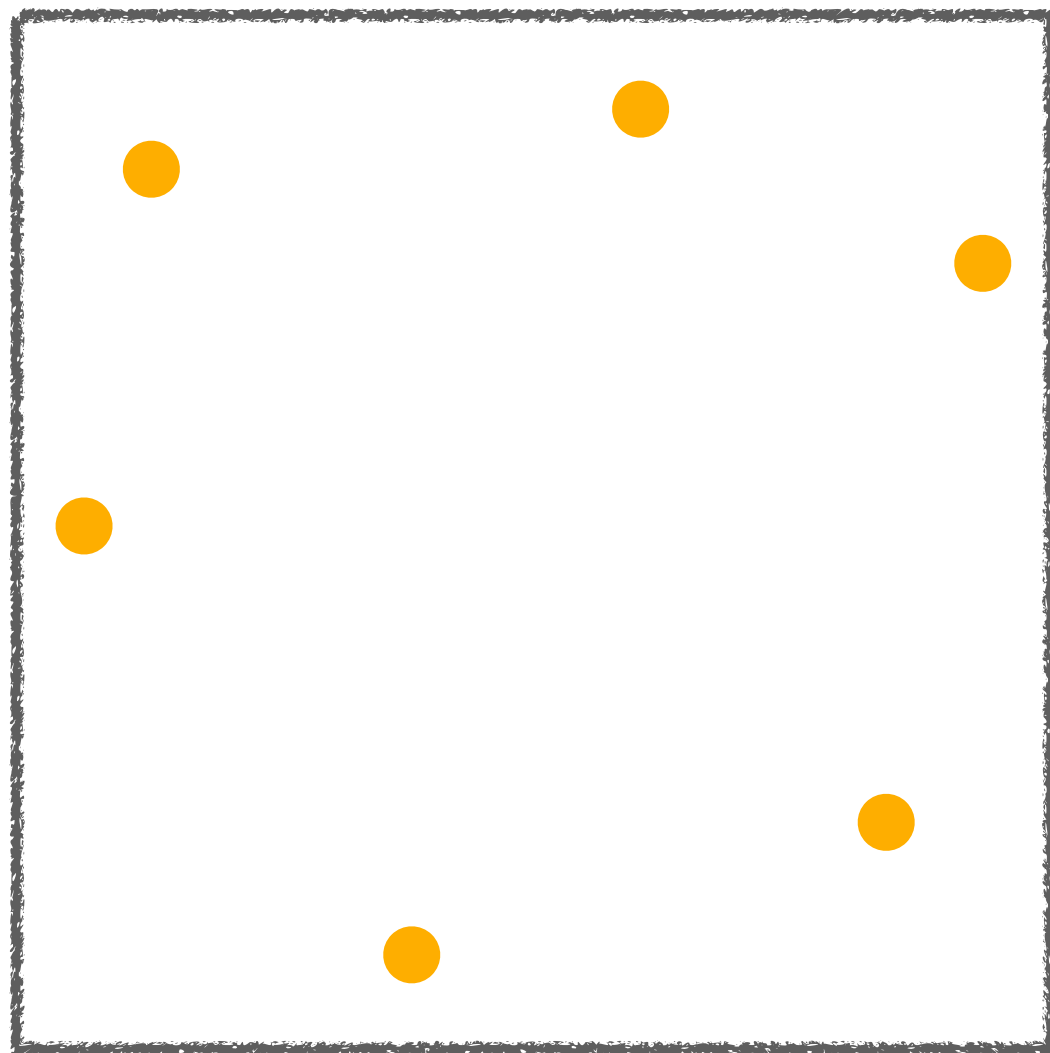
Tutorial #3 — Polygon operations & Farthest point pairs

Peter Kramer

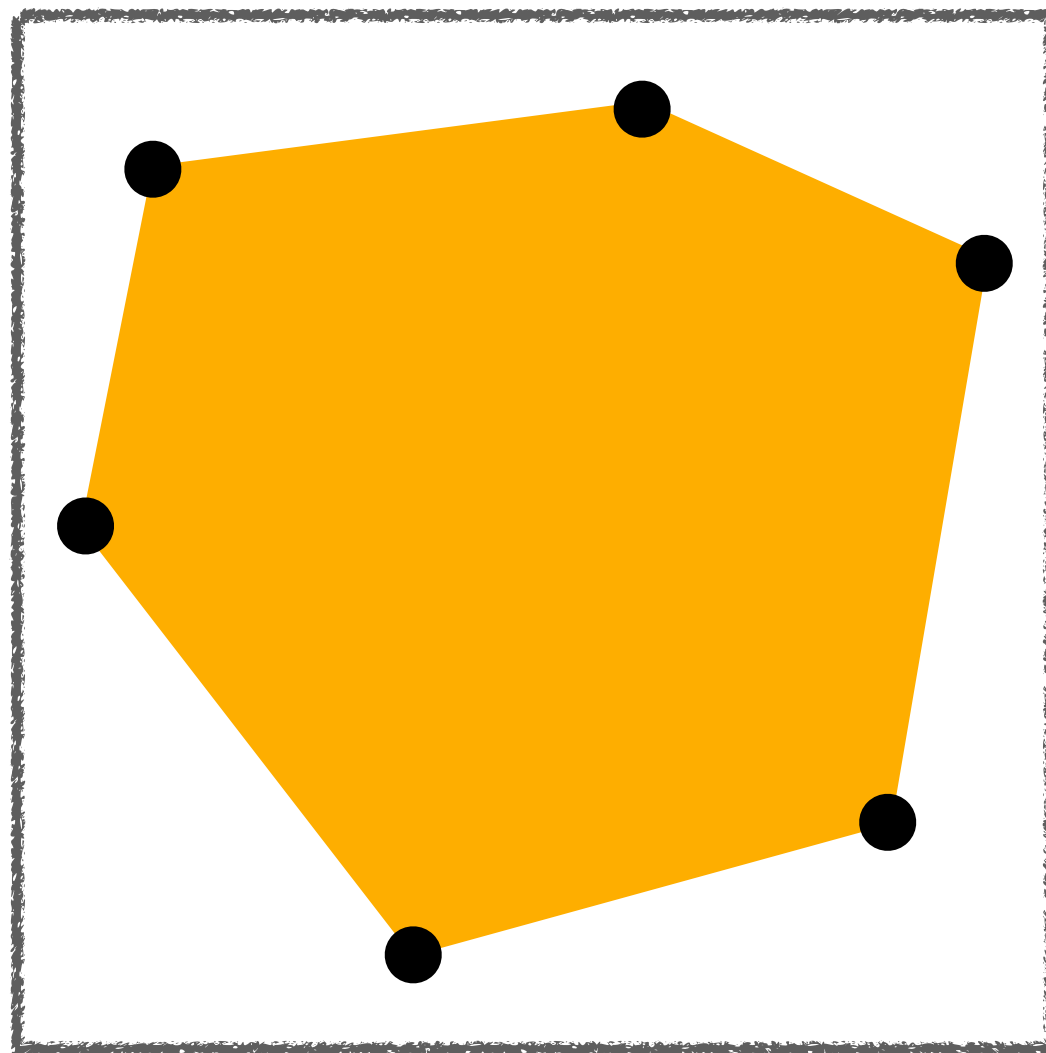
December 7, 2023

Point sets, hulls, and polygons

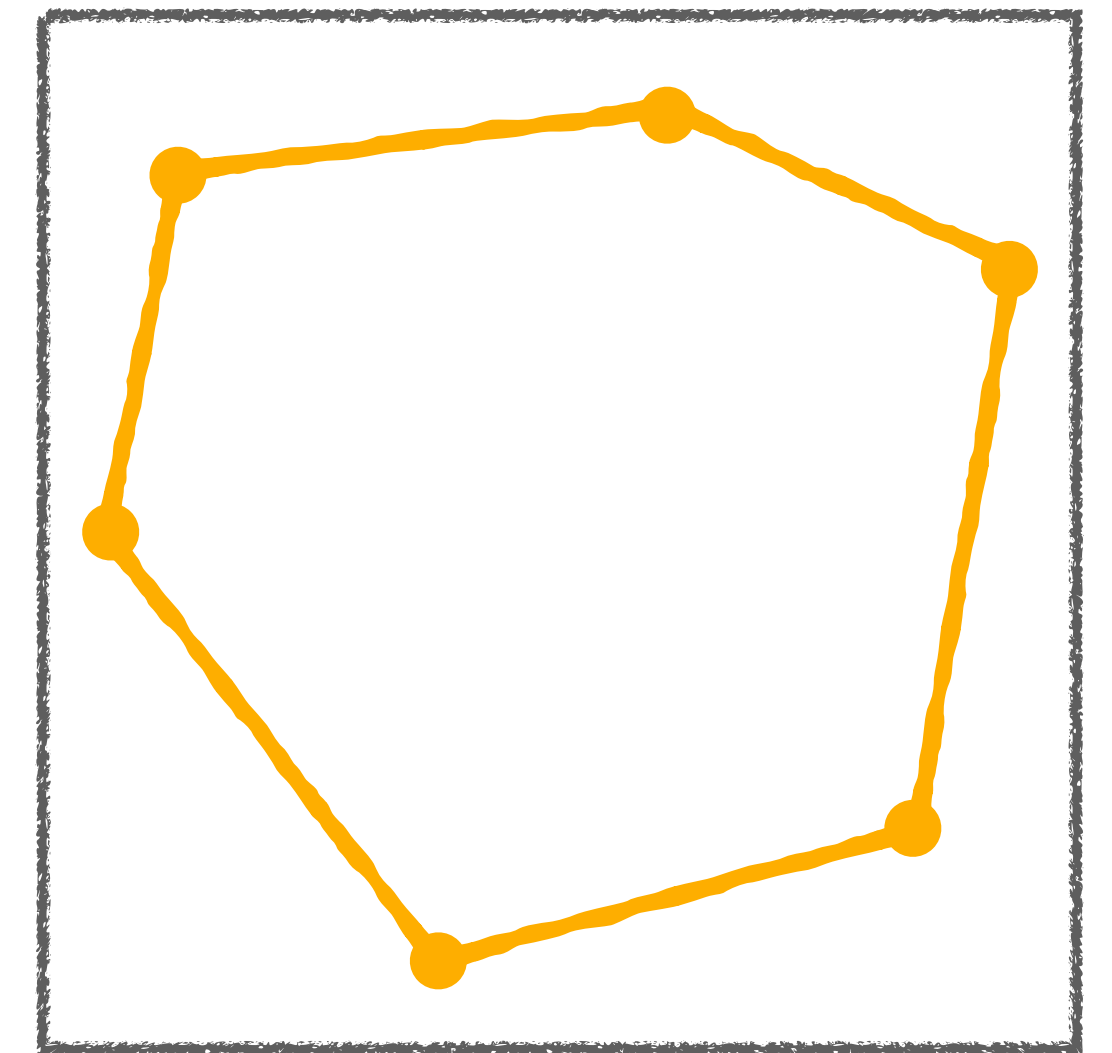
Refresh: What's the difference?



A point set \mathcal{P}



$\text{conv}(\mathcal{P})$



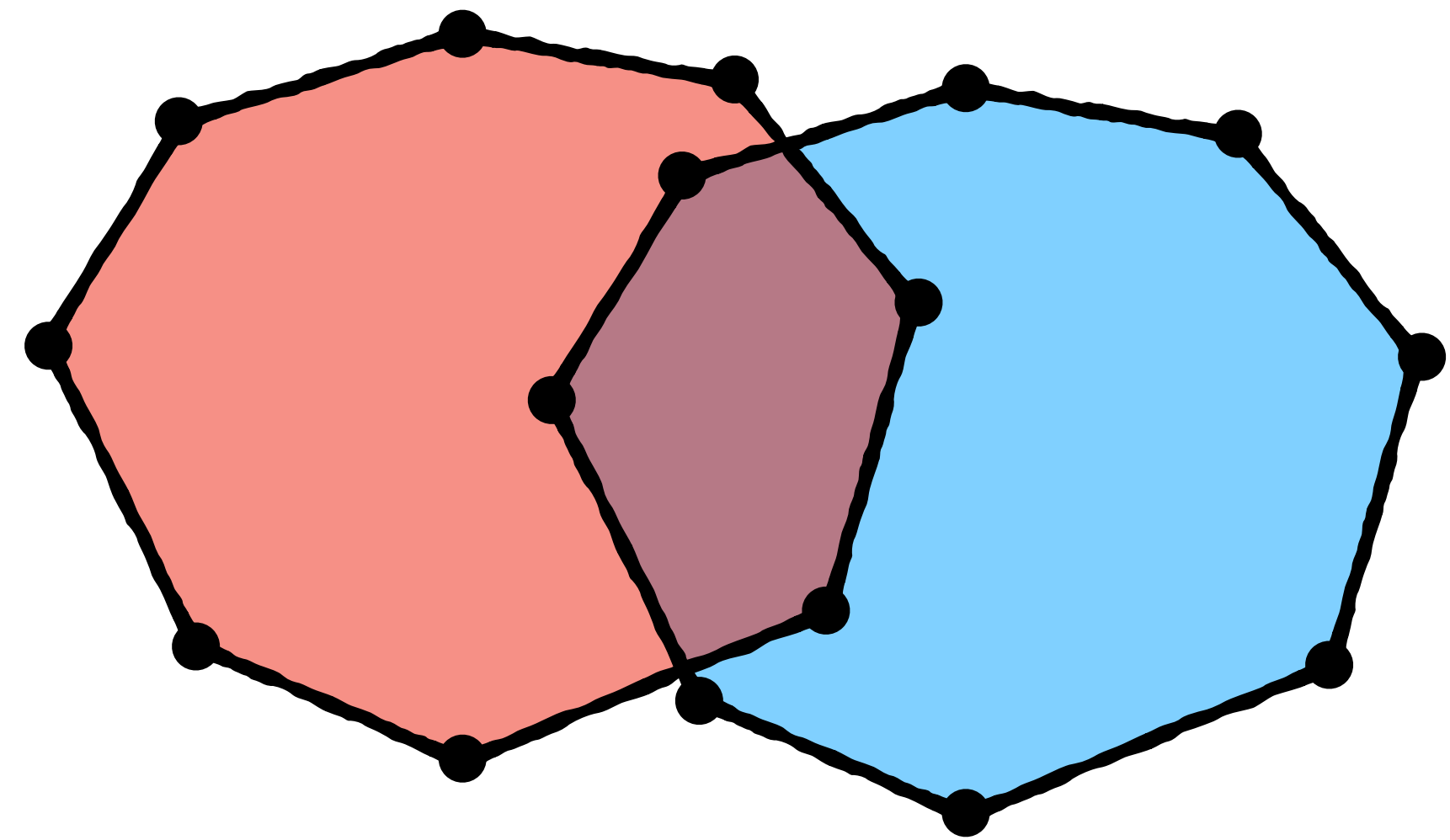
A polygon P on \mathcal{P}

Boolean Operations

Boolean Operations on Convex Polygons

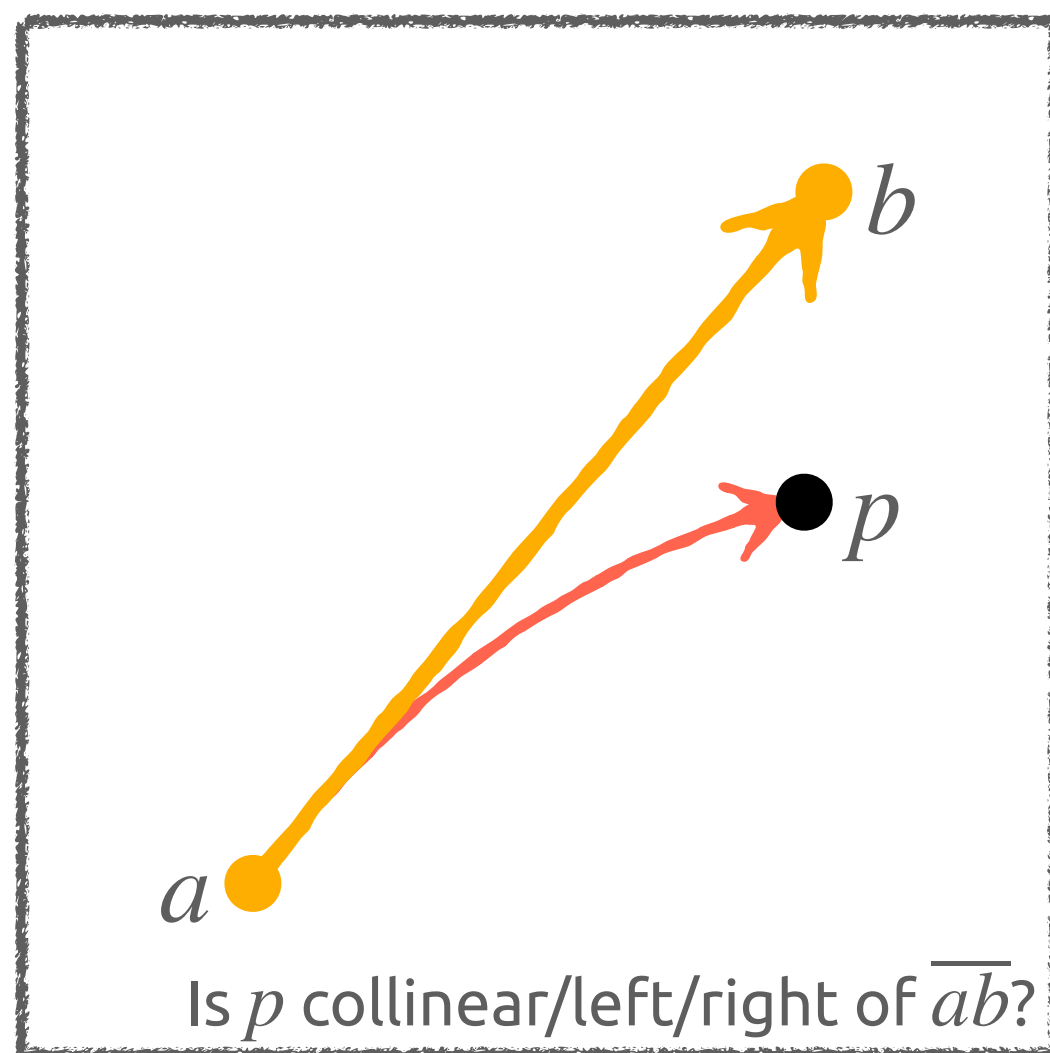
Given two convex Polygons P and Q ,
we seek to determine:

$$P \cap Q, P \cup Q, P \setminus Q, (Q \setminus P)$$



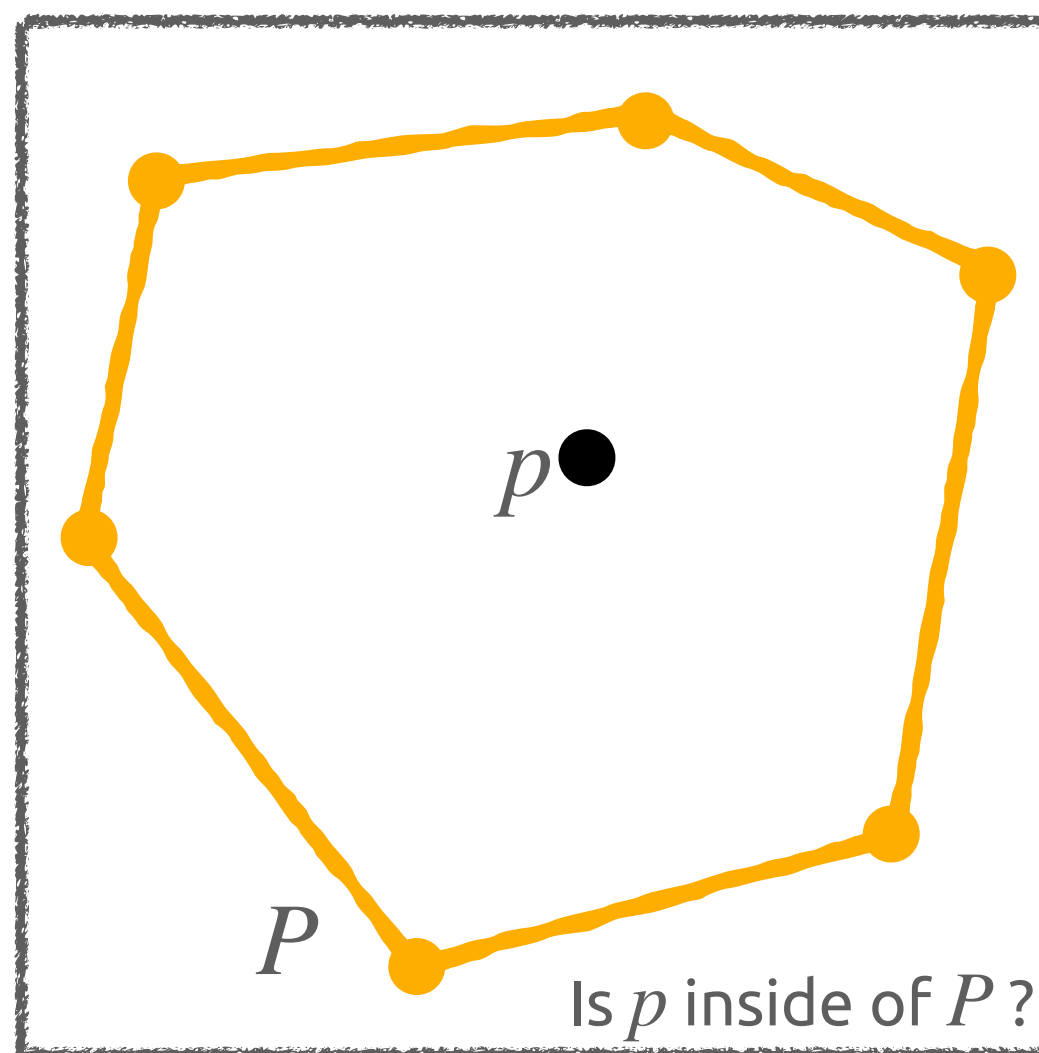
Tools

Point-Line Test



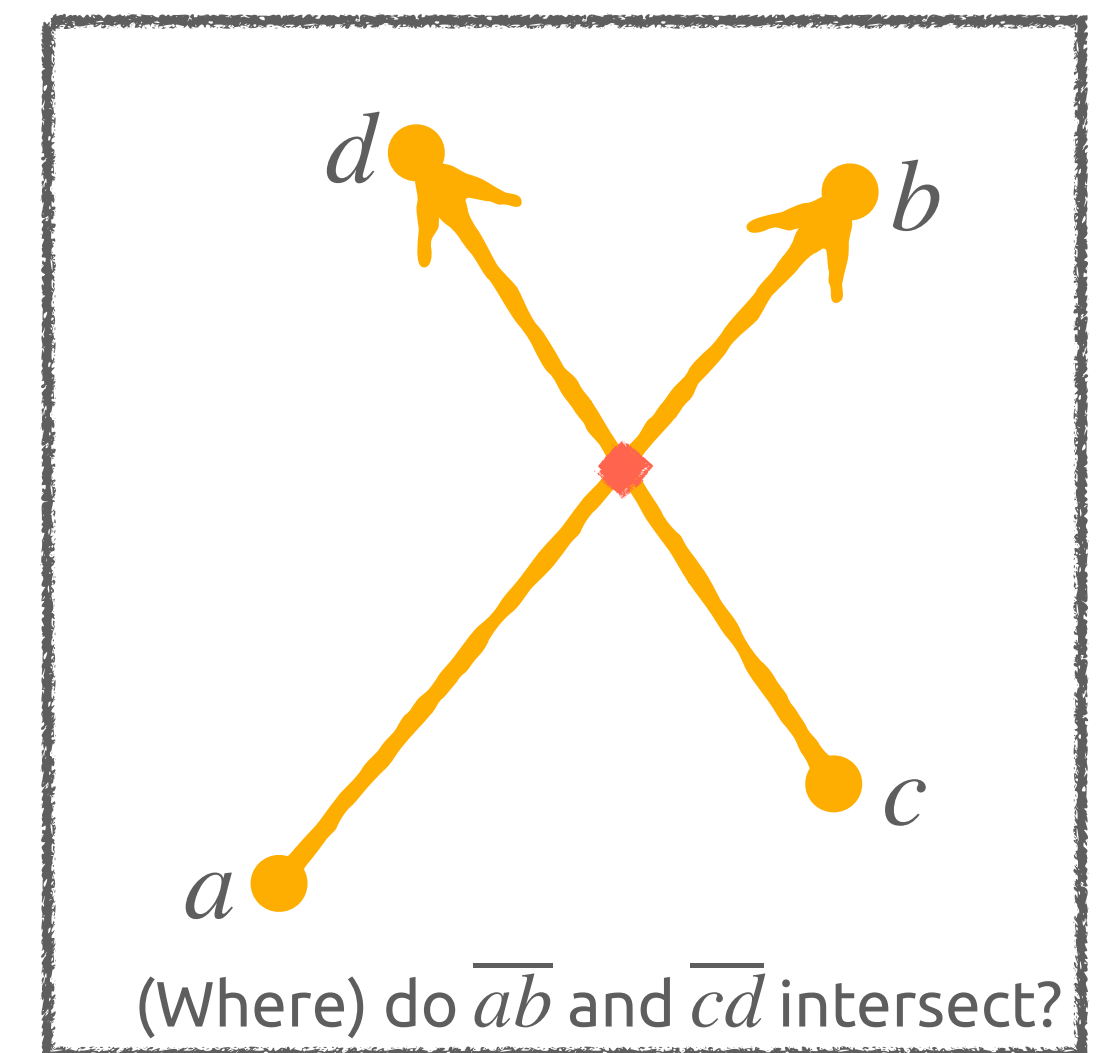
$\mathcal{O}(1)$

Point Location Problem



$\mathcal{O}(\log n)$
(If P convex)

Intersection Test



$\mathcal{O}(1)$

Naive Algorithm using these tools

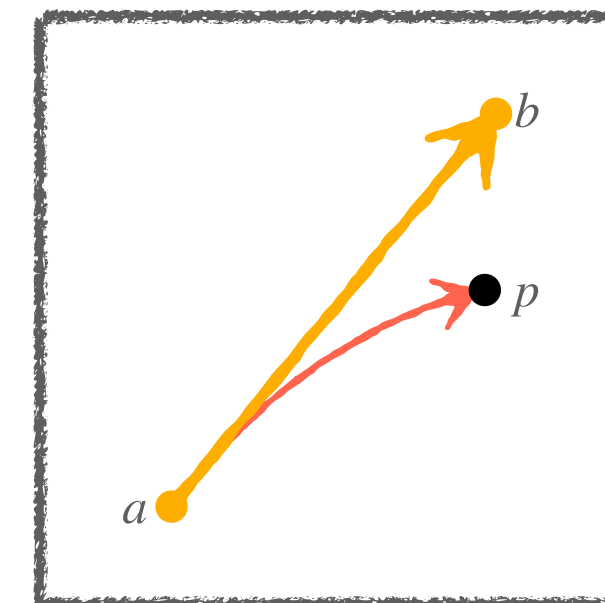
Eliminate/Add points, then recompute convex hull

Given: Convex polygons $P := p_1, \dots, p_n$ and $Q := q_1, \dots, q_m$.

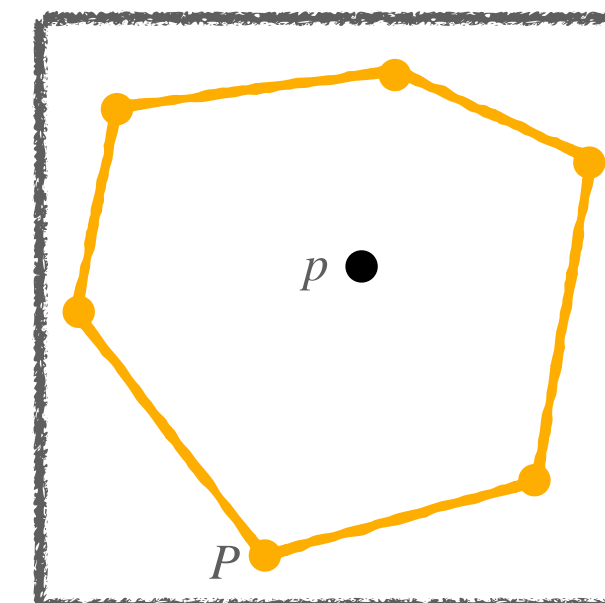
Wanted: The convex polygon $P \cap Q$.

Idea: Determine extreme points of $P \cap Q$ in $\mathcal{O}(n^2)$, then compute the convex hull in $\mathcal{O}(n \log h)$.

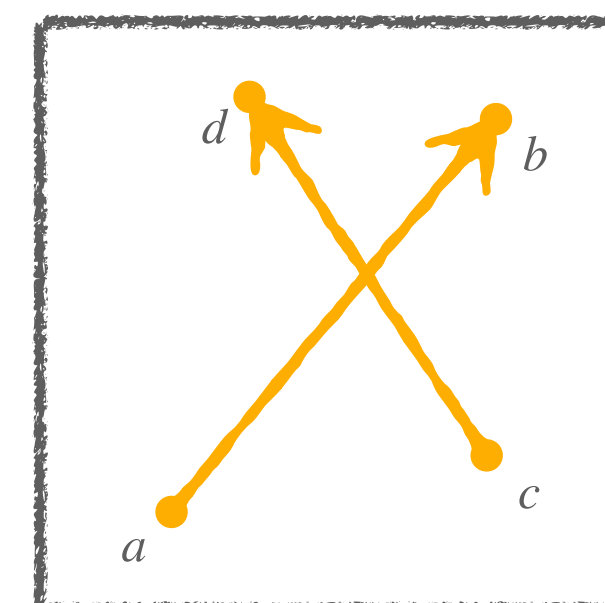
This gives us an $\mathcal{O}(n^2)$ -algorithm.



$\mathcal{O}(1)$



$\mathcal{O}(\log n)$



$\mathcal{O}(1)$

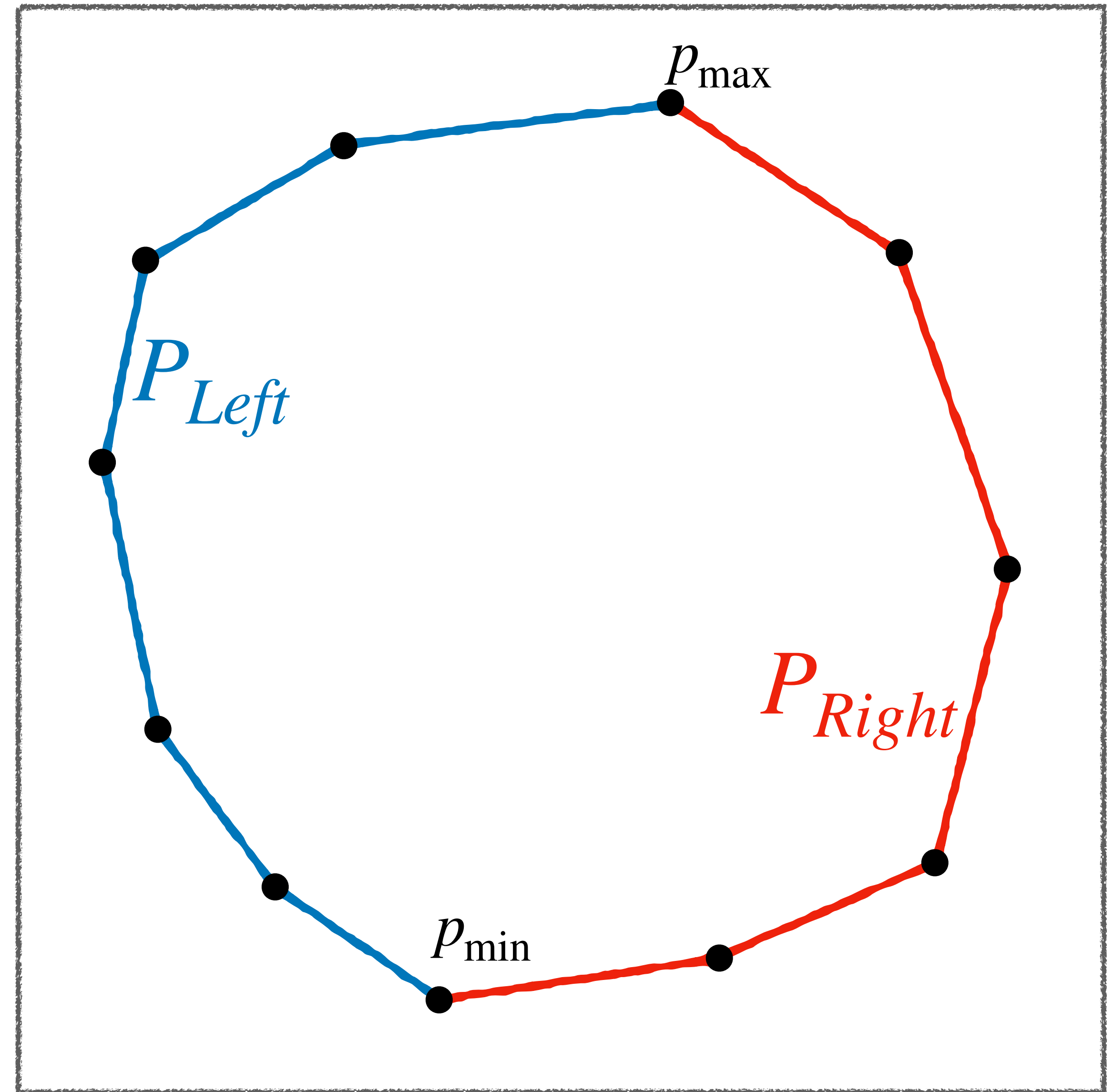
Convex polygons

... can be decomposed!

- Given $P = p_1, \dots, p_n$
- Compute p_{\min} and p_{\max} in $\mathcal{O}(n)$ along the y -axis
- We obtain P_L and P_R :

$$P_L = p_{\max}, \dots, p_{\min}$$

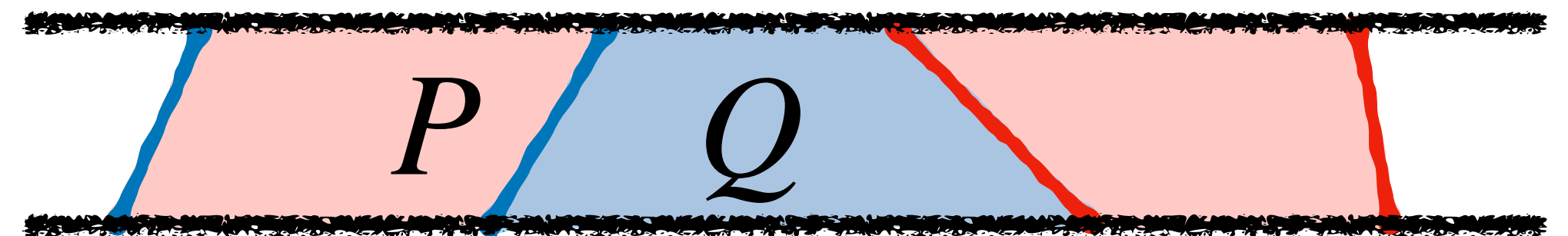
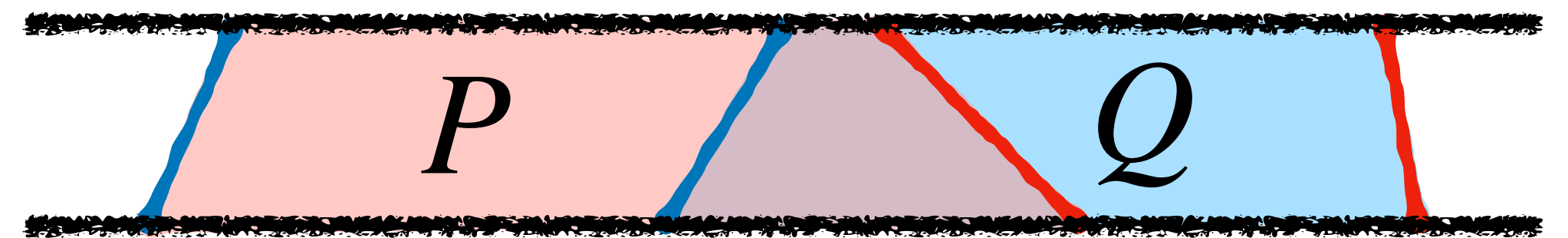
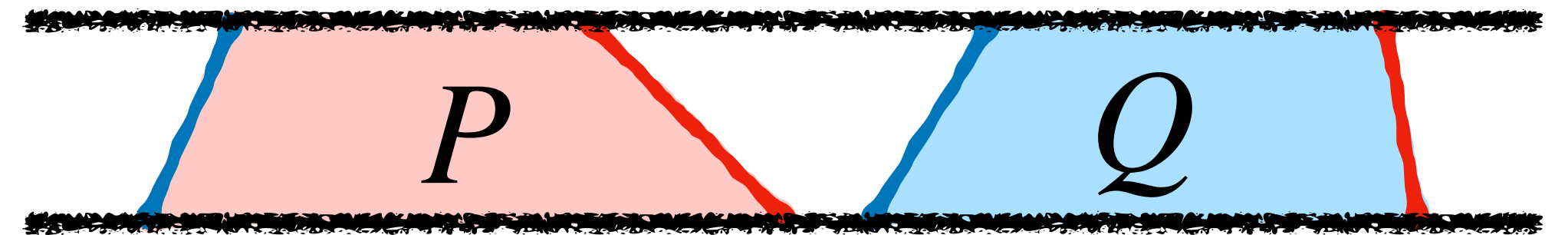
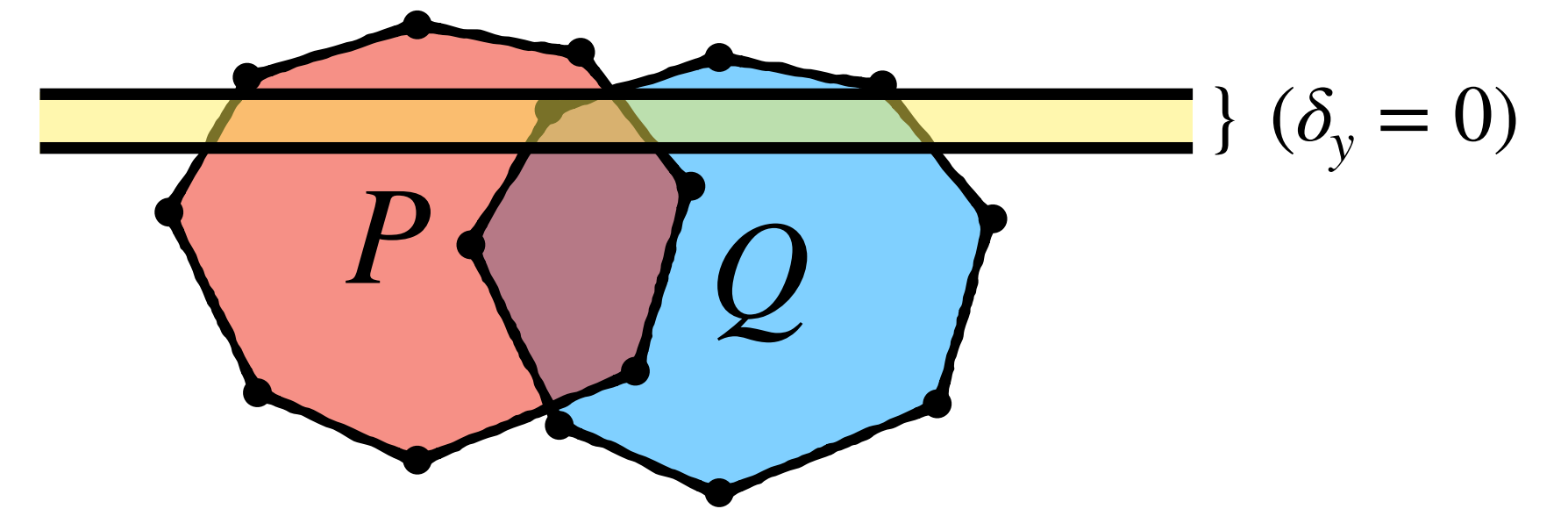
$$P_R = p_{\min}, \dots, p_{\max}$$



Towards an $\mathcal{O}(n)$ -Algorithm

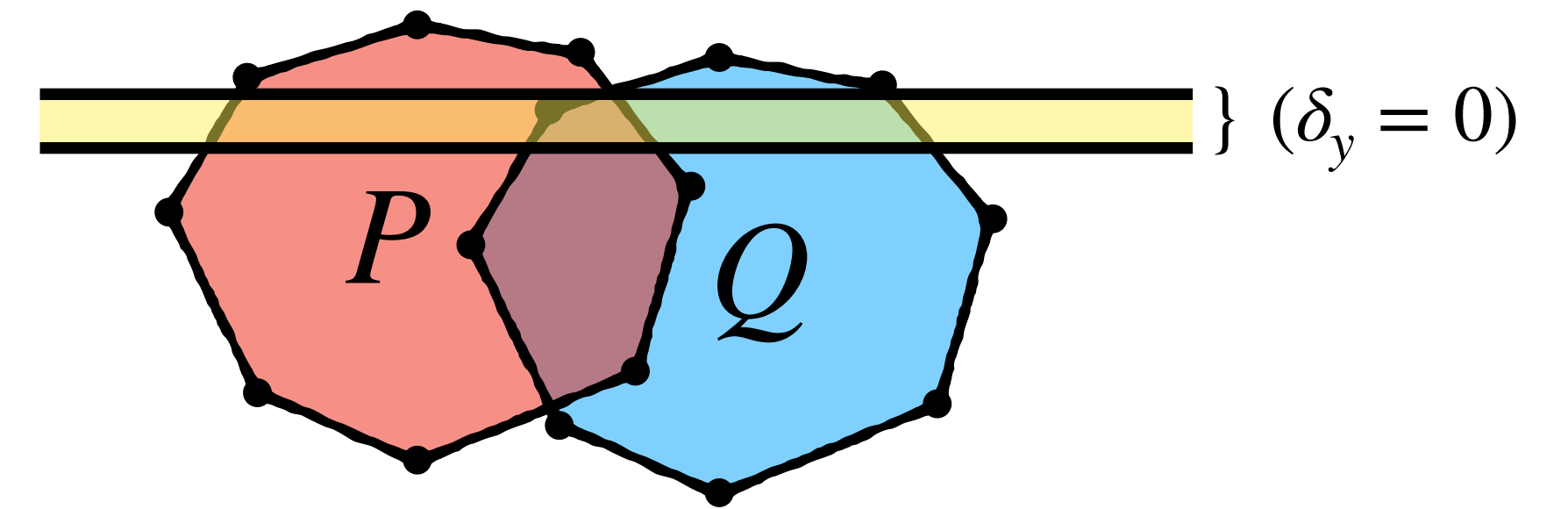
Using left- and right decomposition

- If we slice through P and Q horizontally, at some y (see right, exaggerated):
 - (a) P and Q do not intersect,
 - (b) P and Q overlap partially, or
 - (c) one contains the other.

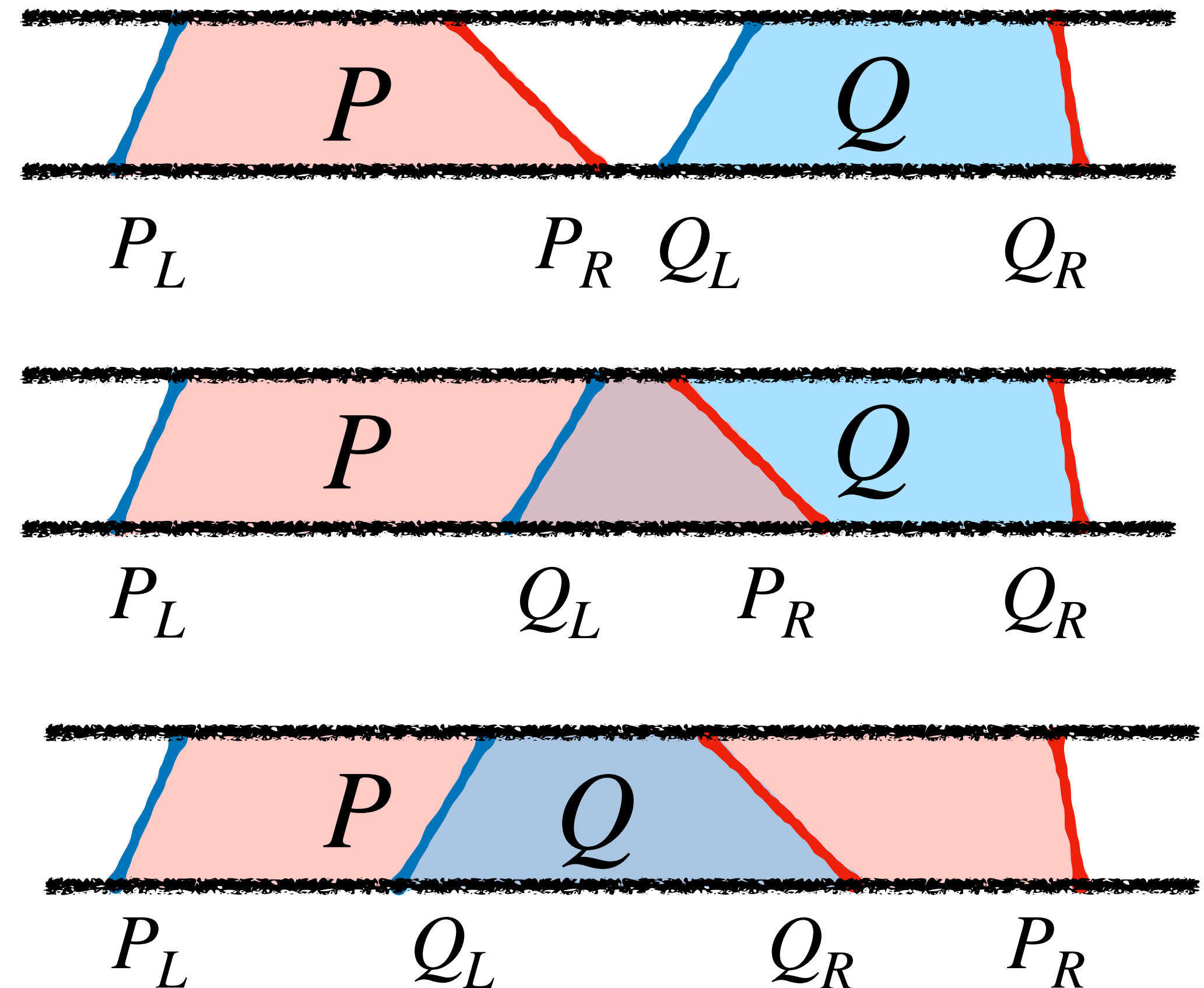


Towards an $\mathcal{O}(n)$ -Algorithm

Using left- and right decomposition



- If we slice through P and Q horizontally, at some y (see right, exaggerated):
 - P and Q do not intersect,
 - P and Q overlap partially, or
 - one contains the other.
- Each case corresponds to an x -order of the chains at that y -coordinate.

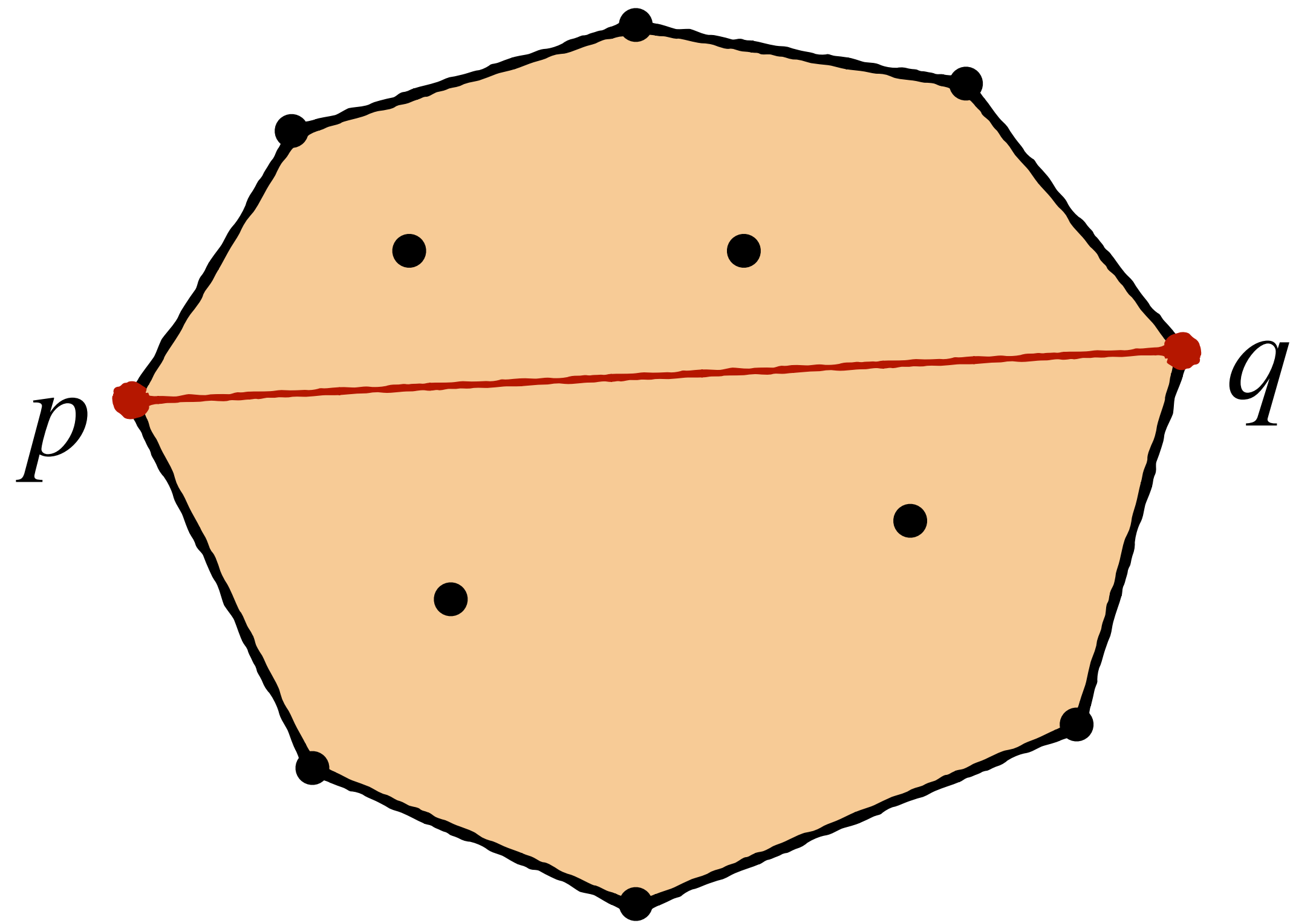


Farthest point pairs

Farthest Point Pairs

Let \mathcal{P} be a finite point set in general position.

The **farthest point pair** of \mathcal{P} consists of two vertices of the convex hull.



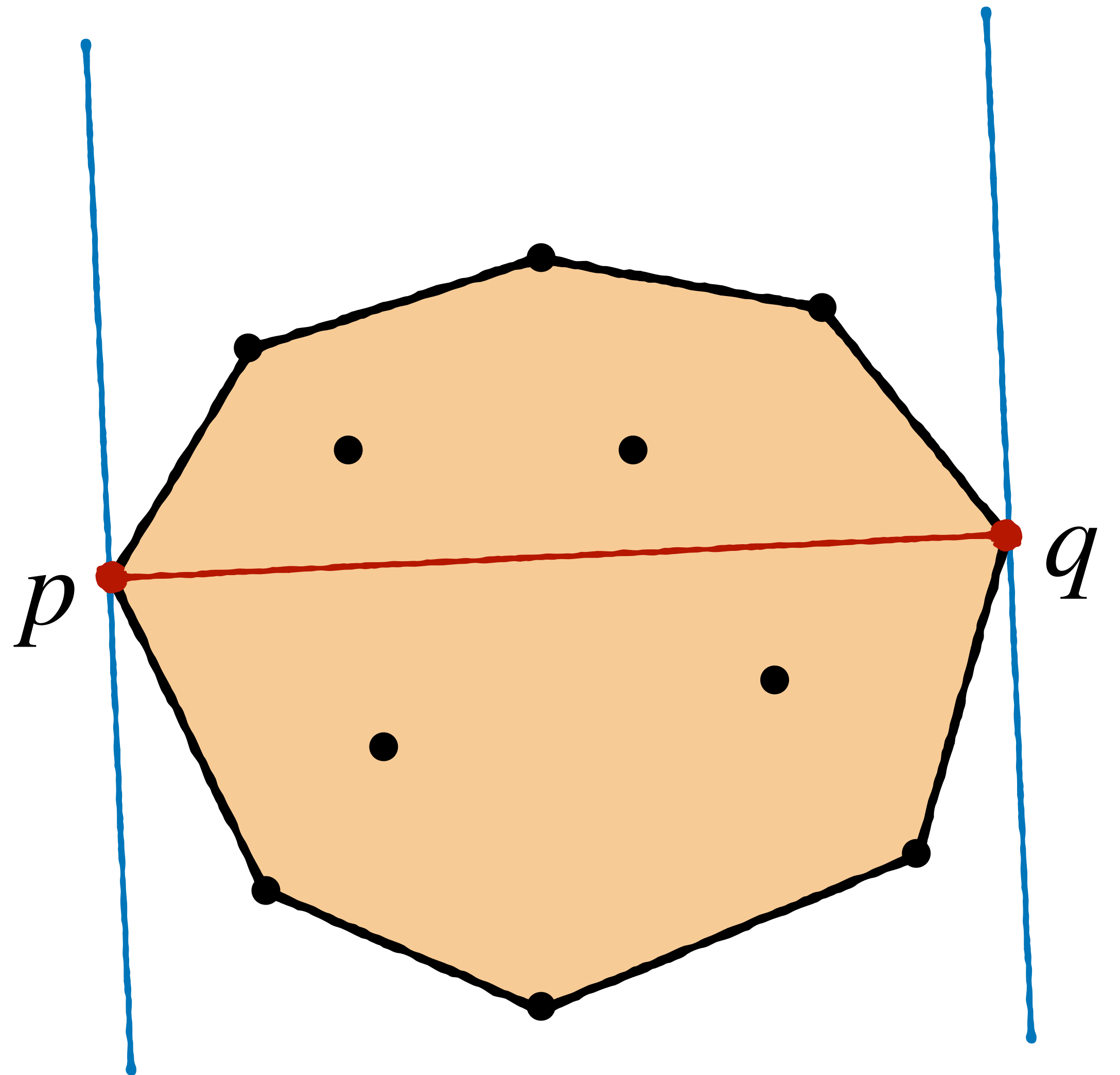
Farthest Point Pairs

Let \mathcal{P} be a finite point set in general position.

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Two points of \mathcal{P} are antipodal if there exist parallel lines through them which do not cut the hull.

Argue that the farthest pair is antipodal.



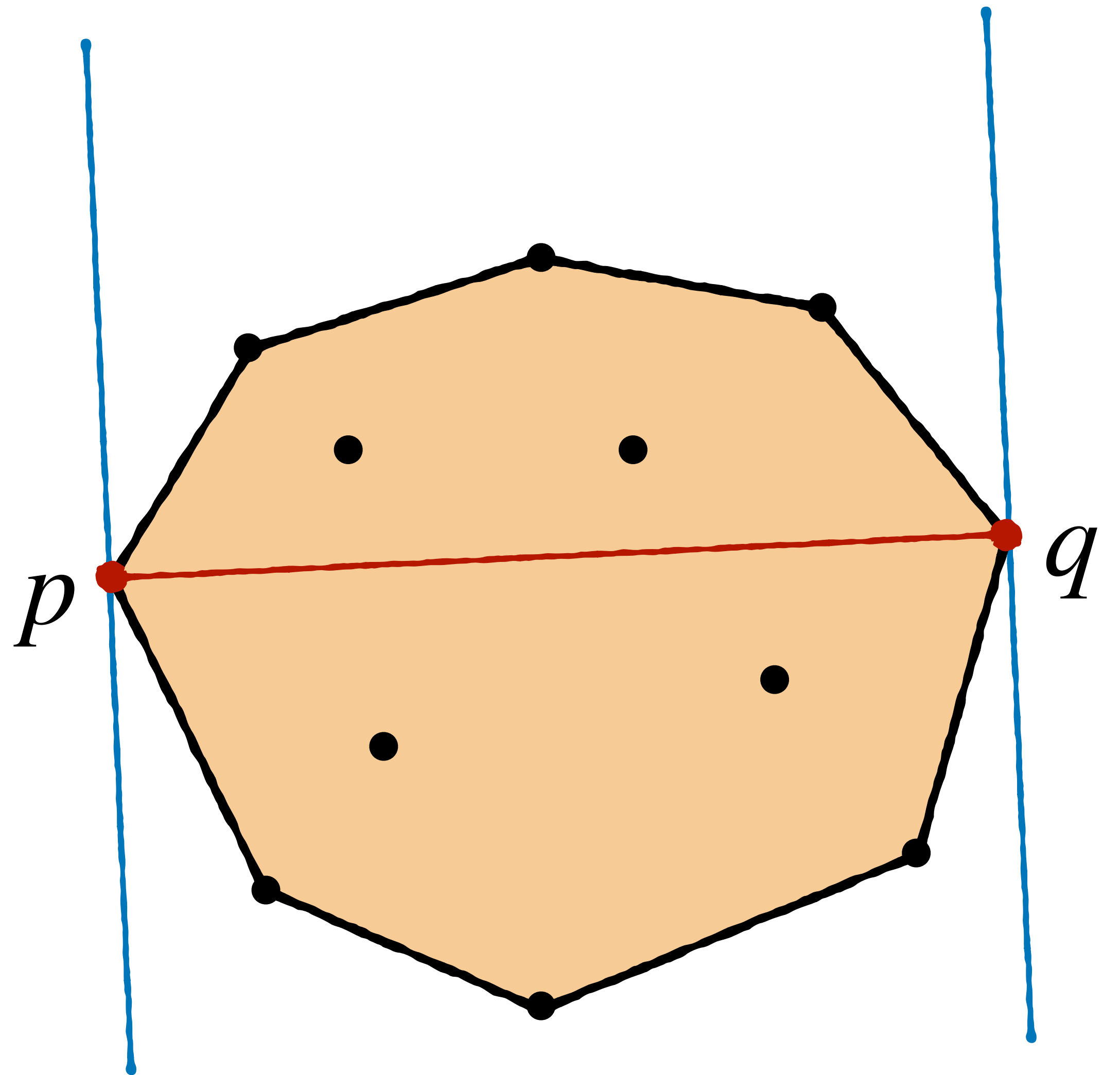
Farthest Point Pairs

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How can we use this?

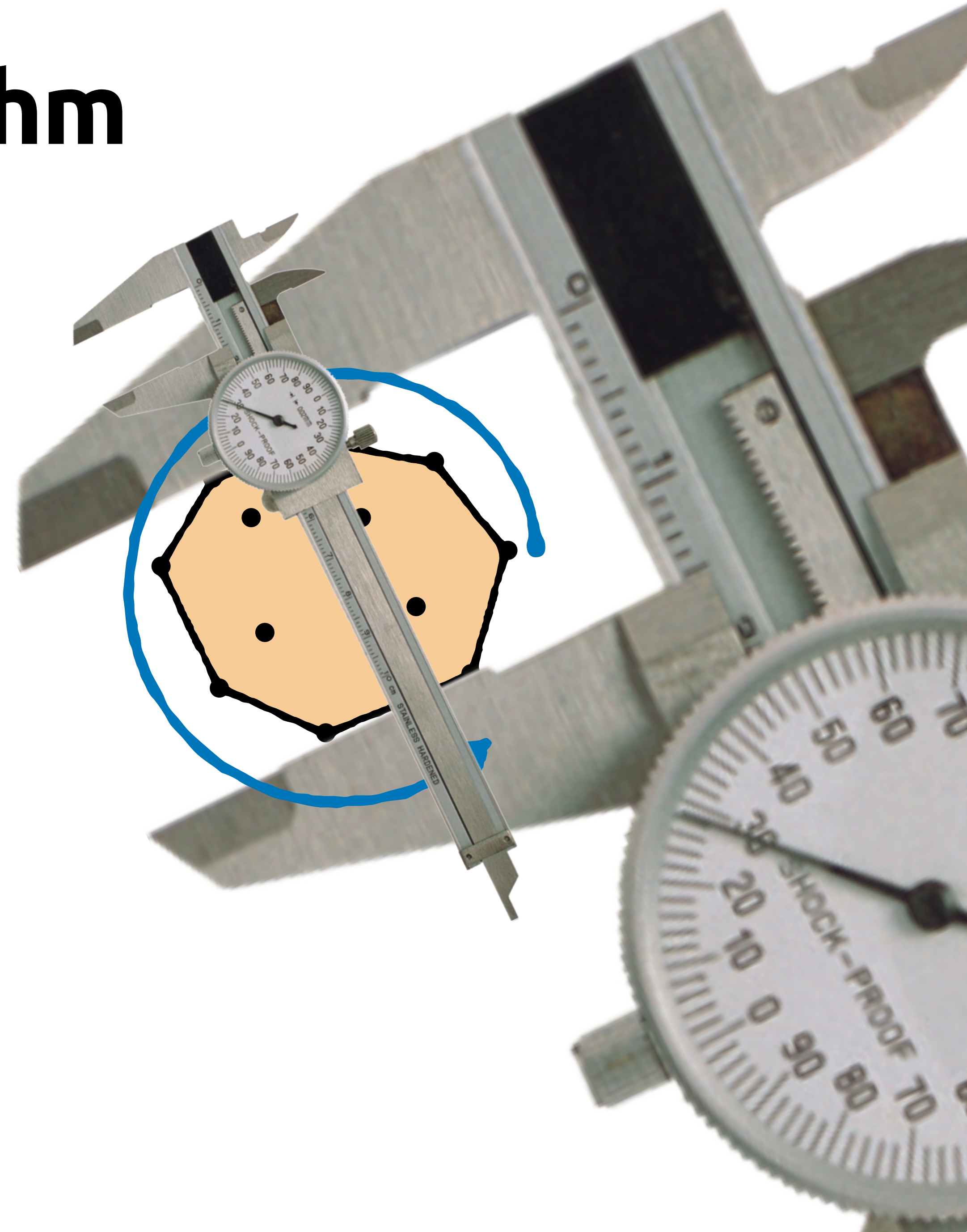


Rotating Callipers Algorithm

Michael Shamos, 1978

Idea: Compute convex hull, then enumerate antipodal pairs and track the farthest one.

To achieve this, “rotate” parallel lines around the point set.



Rotating Callipers Algorithm

Michael Shamos, 1978

Distances [\[edit\]](#)

- Diameter (maximum width) of a convex polygon^{[6][7]}
- Maximum distance between two convex polygons^{[9][10]}
- **Minimum distance** between two convex polygons^{[11][12]}

machine

Bounding boxes [\[edit\]](#)

- Minimum area [oriented bounding box](#)
- Minimum perimeter [oriented bounding box](#)

Triangulations [\[edit\]](#)

- Onion [triangulations](#)
- Spiral [triangulations](#)
- [Quadrangulation](#)
- Nice triangulation
- Art gallery problem
- Wedge placement optimization problem^[15]

Multi-polygon operations [\[edit\]](#)

- Union of two convex polygons
- Common tangents to two convex polygons
- Intersection of two convex polygons^[16]
- [Critical support lines](#) of two convex polygons
- Vector sums (or Minkowski sum) of two convex polygons^[17]
- Convex hull of two convex polygons

Traversals [\[edit\]](#)

- Shortest transversals^{[18][19]}
- Thinnest-strip transversals^[20]

Others [\[edit\]](#)

- Non parametric decision rules for machine learned classification^[21]
- Aperture angle optimizations for visibility problems in computer vision^[22]
- Finding longest cells in millions of biological cells^[23]
- Comparing precision of two people at firing range
- Classify sections of brain from scan images

