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RAIM: Redundant Array of Independent Motes

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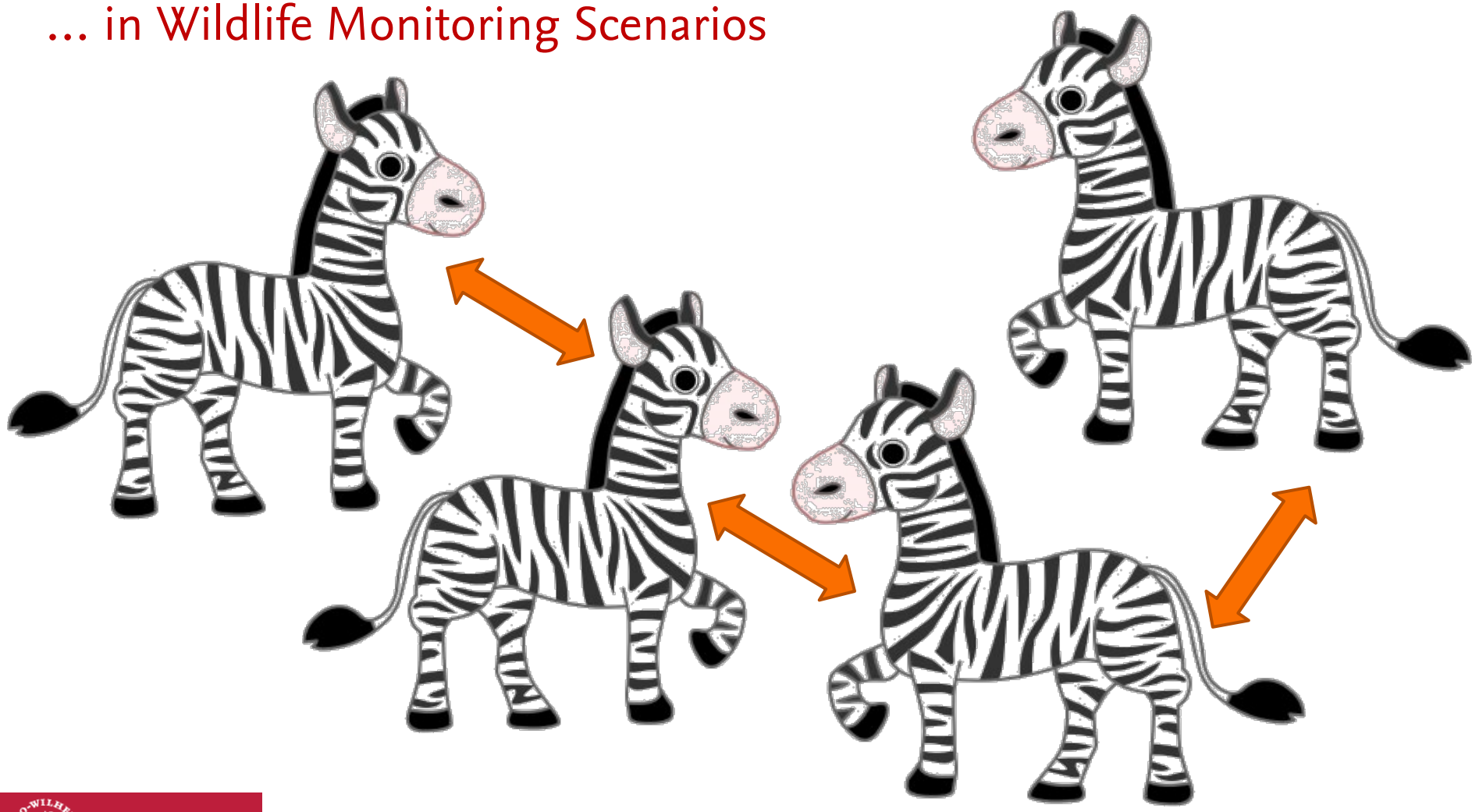
Motivation: Store Data on WSN nodes

… in Wildlife Monitoring Scenarios

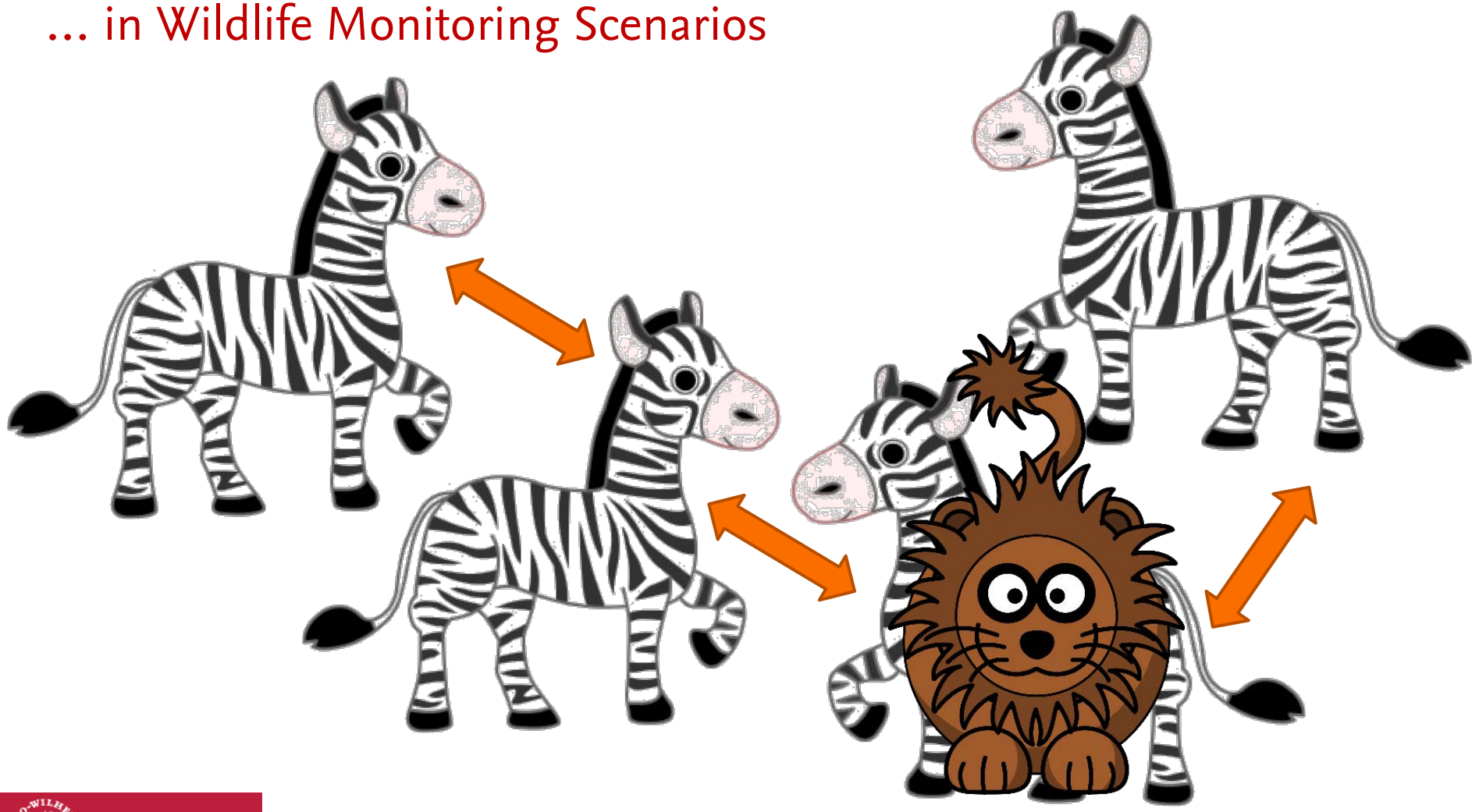


Wild LION

appeared!

Motivation: Store Data on WSN nodes

... in Wildlife Monitoring Scenarios



Motivation: Store Data on WSN nodes

... in Personal Area Networks

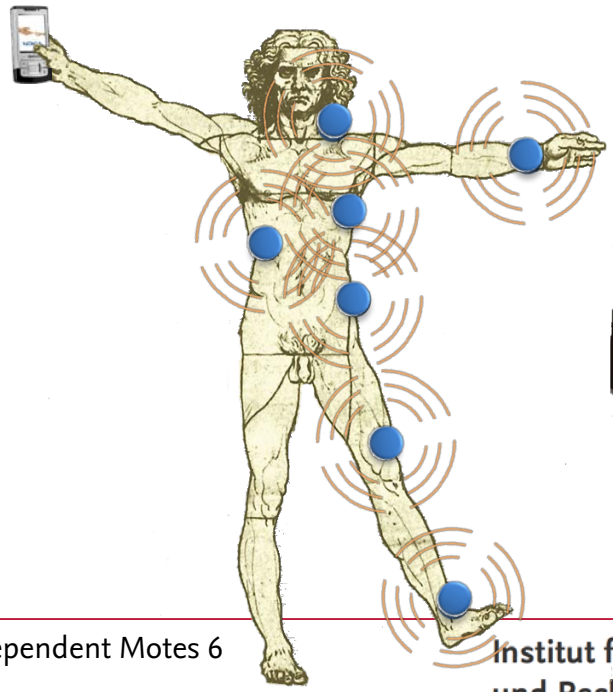
WSN Challenges

Nodes can get

- Lost
- Stolen
- Eaten
- Destroyed

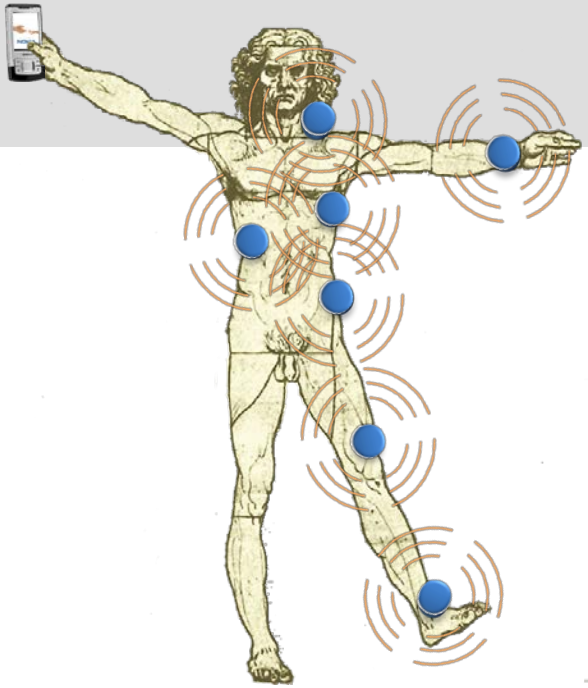
Data can be

- Private / Confidential



- Data should be preserved

So, what's the question here?



How can we ensure

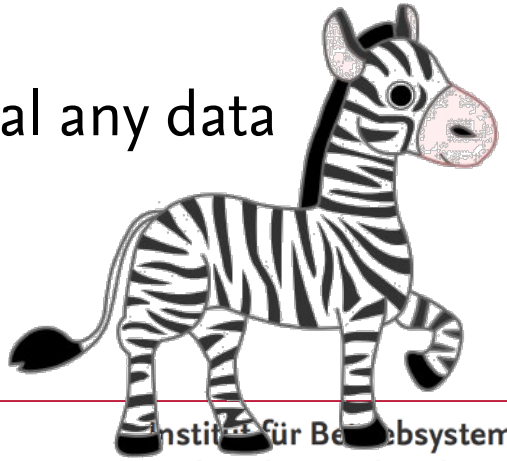
- Privacy and Confidentiality
- Data Security

… at the same time

… and in different configurations?

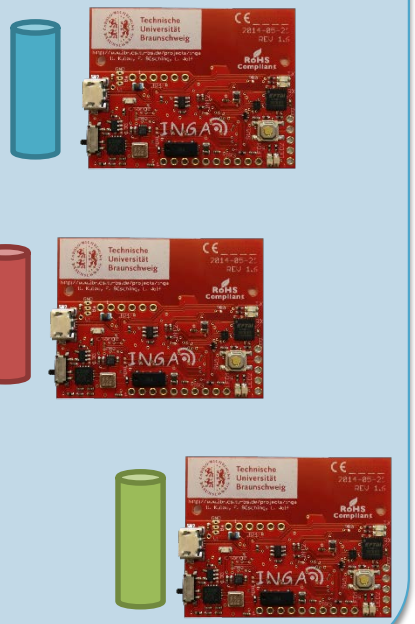
Can we achieve different levels of Redundancy and Confidentiality?!

- A single "missing" BAN sensor should not reveal any data
- Collecting one Zebra is sufficient to get all data
- → depending on the individual scenario!



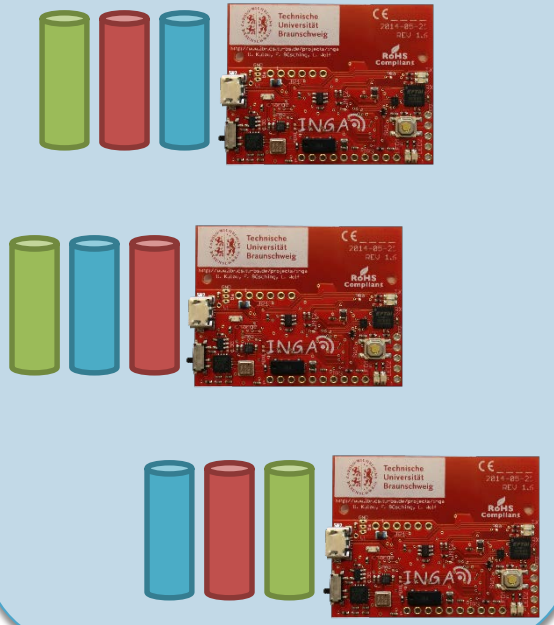
Idea

Individual data stored on individual nodes



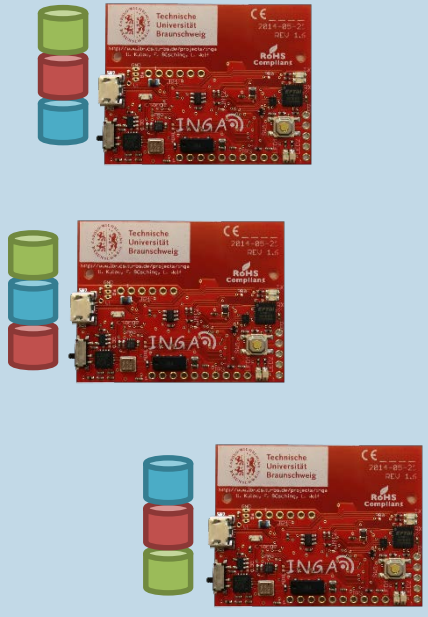
No redundancy

All data stored on every node



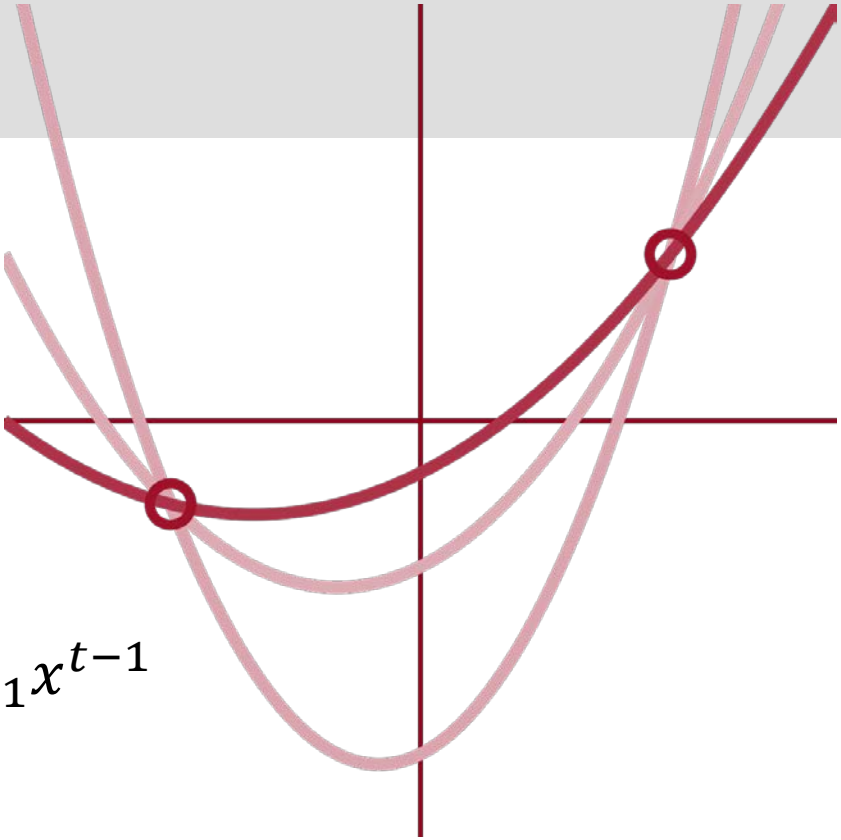
Full redundancy

Parts of data stored on each node



Full confidentiality

Basics: Secret Sharing



„normal“ Secret Sharing

- $s = s_1 + s_2 + \cdots + s_n$
- All parts of s needed

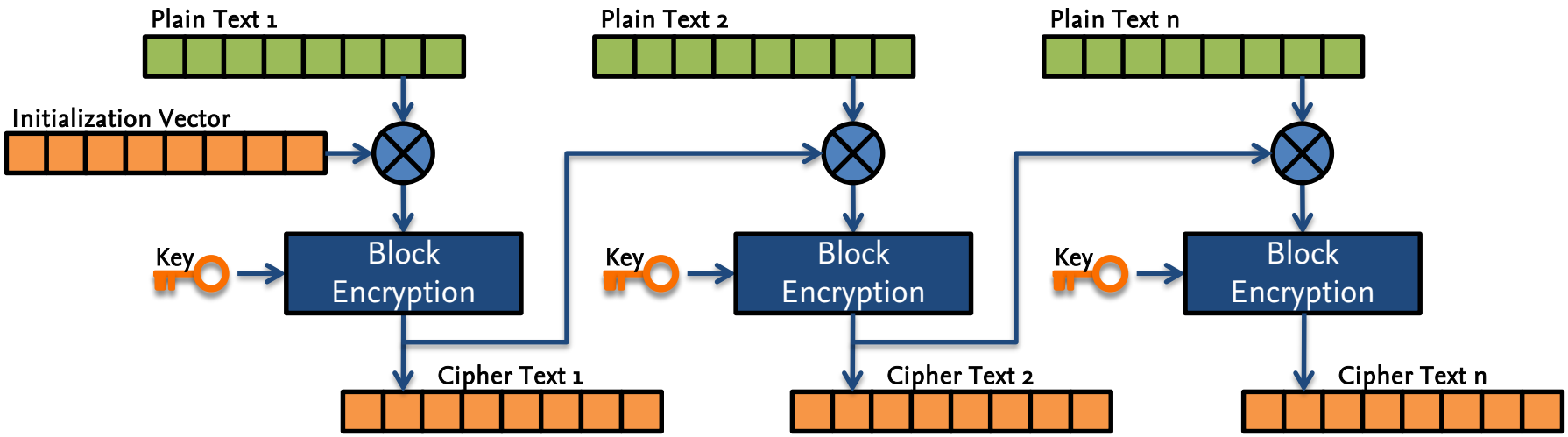
Shamir’s Secret Sharing

- $f(x) = s + a_1 x + a_2 x^2 + \cdots + a_{t-1} x^{t-1}$
- t pairs of values $\left(x, f(x)\right)$ needed
- $t - 1$ keys can be compromised

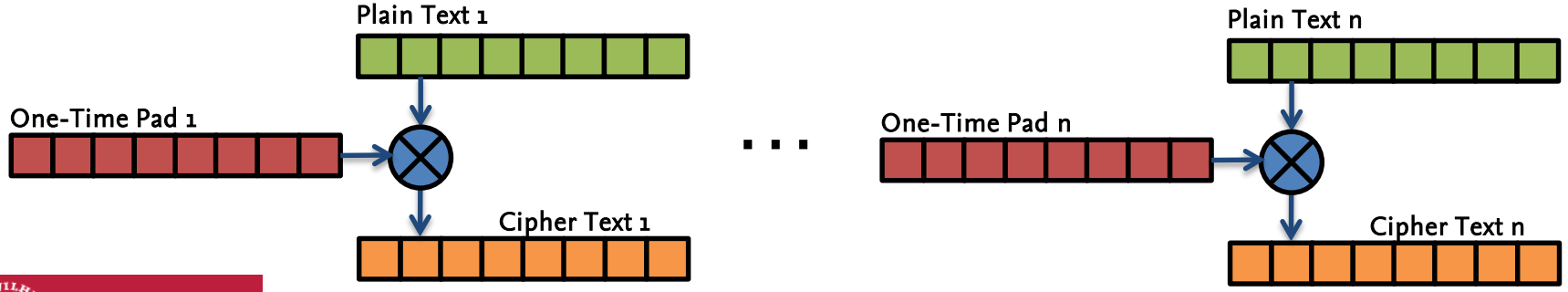
Optimization for µC

- $f(x) = s + a_1 x + a_2 x^2 + \cdots + a_{t-1} x^{t-1} \; mod \; p$

AES in CBC mode (Chiper Block Chaining)



One-Time Pads



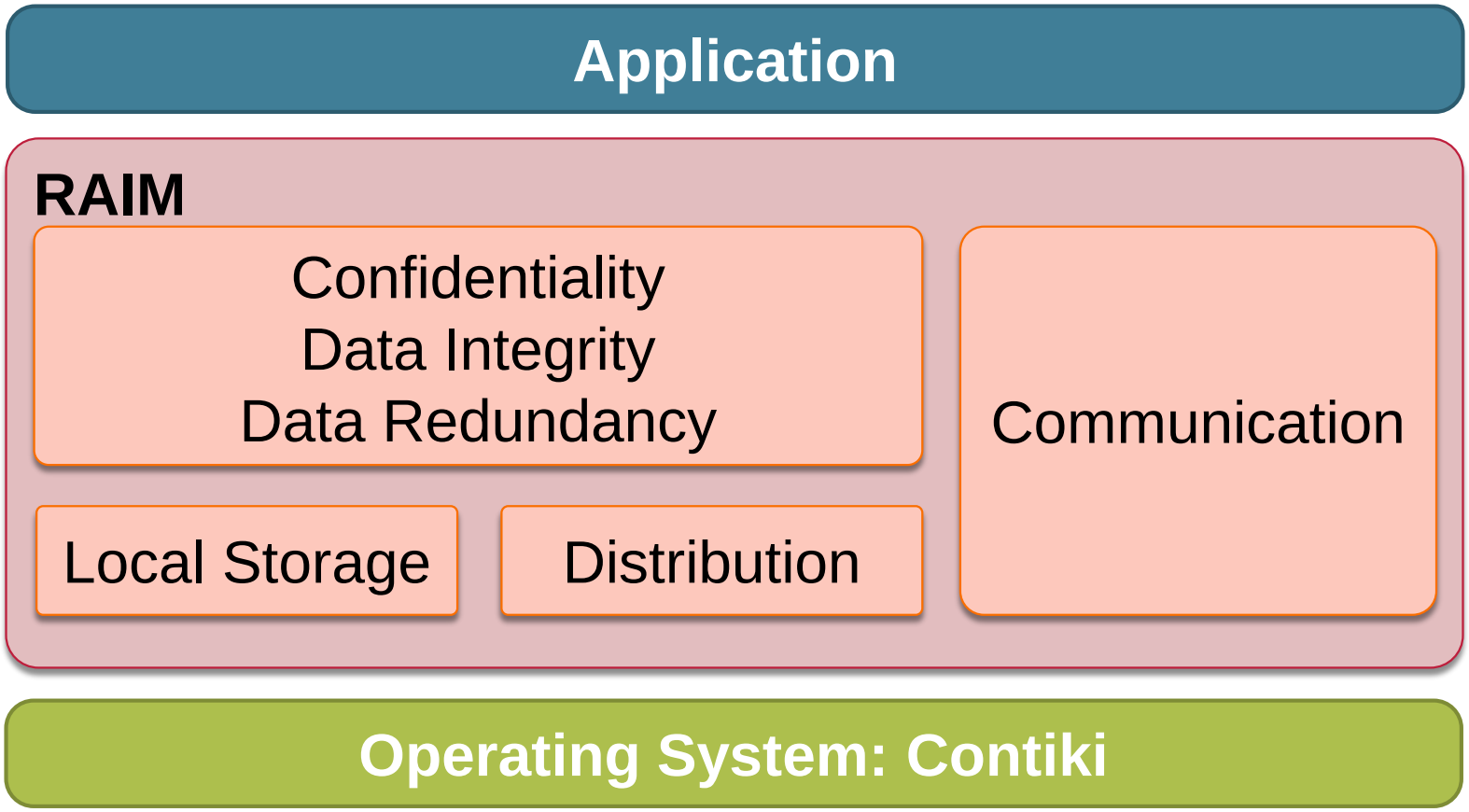
So, what did we do?

Take

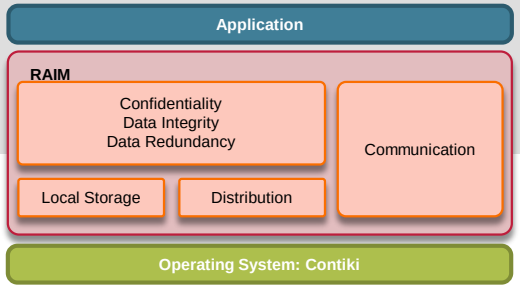
- Shamir‘s Secret Sharing
- Well known encryption methods
- Well known authentication methods

... and build it.

RAIM's modular architecture



Implementation – I



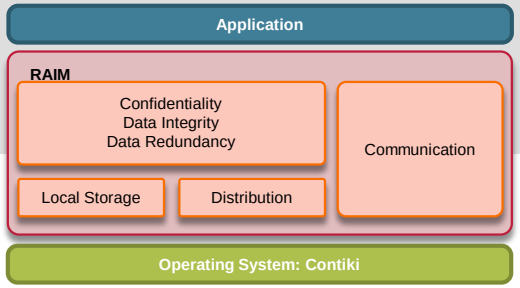
Operating System

- Contiki

Local storage

- SD-card
- FAT 32
- Special directory structure:
  - *local* – contains locally generated data
    - In OTP configuration also OTPs are stored here
  - *$mote_id* – contains data received from other nodes

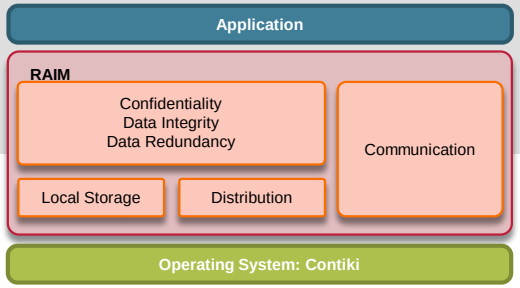
Implementation – II



Confidentiality

- AES
  - In CBC-mode
  - Hardware AES in RF233
  - Pre-shared key
- One-Time Pads
  - In place
- OTP needs large storage
  - But not additional
- AES can be used for nodes without storage

Implementation – III

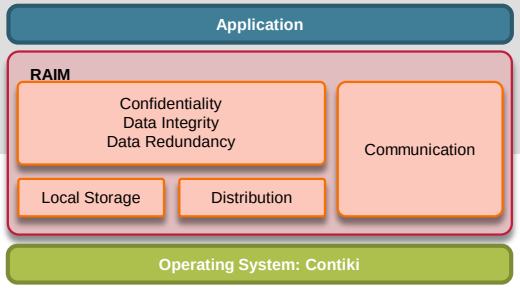


Integrity

- CBC-MAC
  - Hardware AES in RF233
  - Tag appended to transmitted data

- CRC-8 checksum
  - When no pre-shared key is available
  - Only error detection

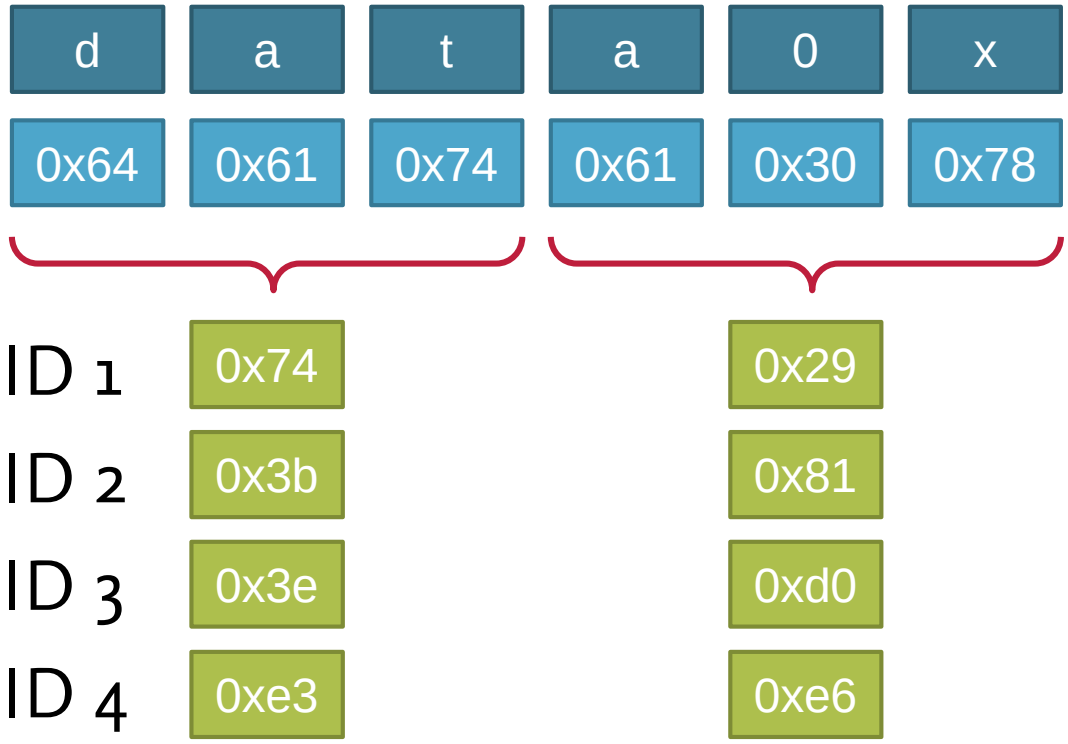
Implementation – IV



Data Redundancy and Distribution

- No redundancy
  - Data is only stored (and encrypted) locally
- Full redundancy
  - All data is spread to any other nodes
- (k;n)-threshold Sharmir's Secret Sharing
  - Data is spread over n nodes
  - k nodes are needed for encryption

Examplary Distribution: (3;4)-threshold SSS scheme

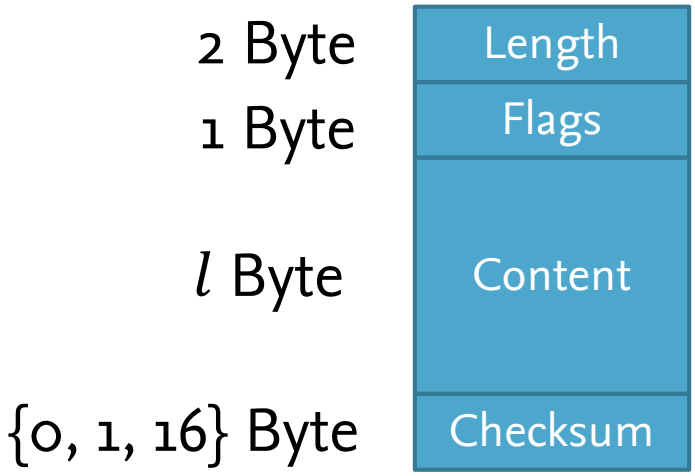


- 6 Byte divided in
- 2 chunks a 3 Byte

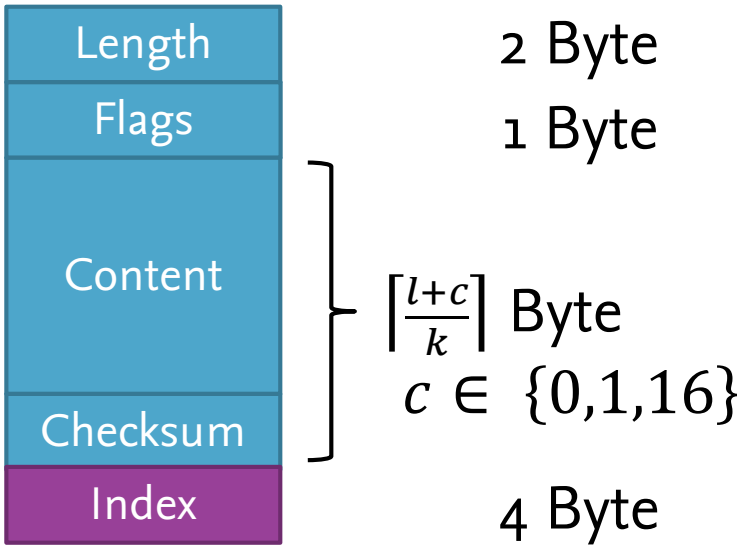
- chunks distributed among 4 motes

3 out of 4 motes are required to reconstruct data

Evaluation: Storage Requirements For l Byte of Data



Local storage of sensing mote



Storage of neighboring mote (remote)

Cipher Suite Configurations

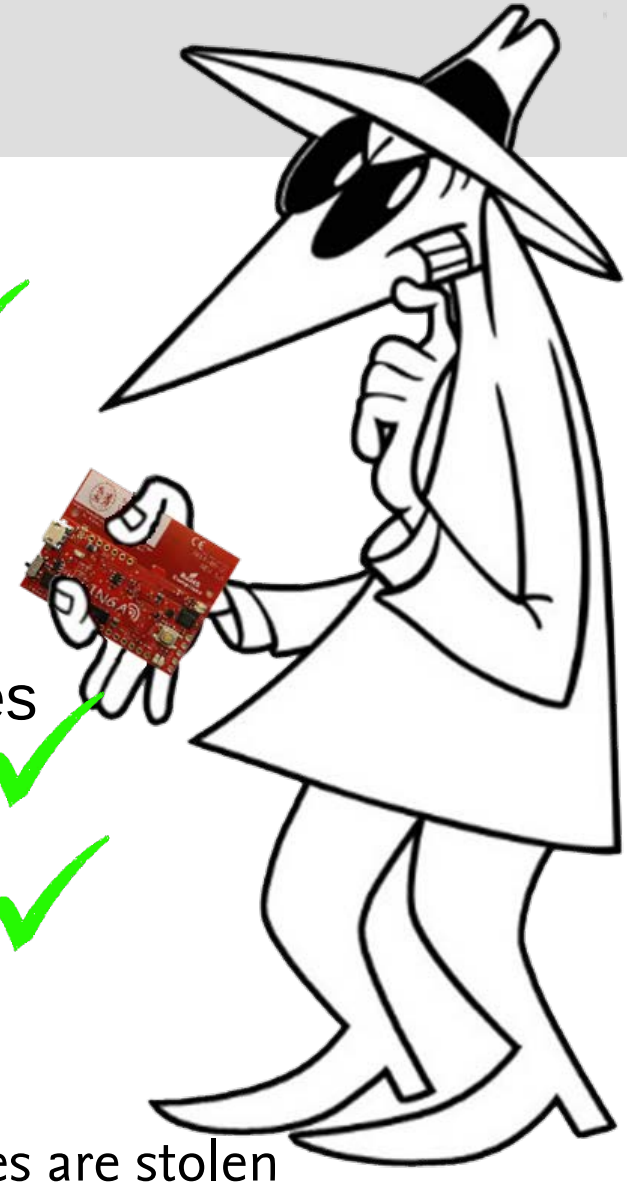
Cipher Suite	Properties
CRC only	Error Detection
CBC-MAC only	Data Integrity
AES	Confidentiality
OTP	Confidentiality
OTP + CBC-MAC	Confidentiality + Data Integrity
AES + CBC-MAC	Confidentiality + Data Integrity

Attack Models - External Attacker

- Eavesdropping/man-in-the-middle
  - If configured for OTP/AES + CBC-MAC
  - → Data encrypted and authenticated ✓
- Insert data
  - If configured for OTP/AES
  - → Preshared key is not known ✓
- Replay attacks
  - Index included and authenticated (CBC-MAC)
  - →receiving nodes verify index not used before ✓

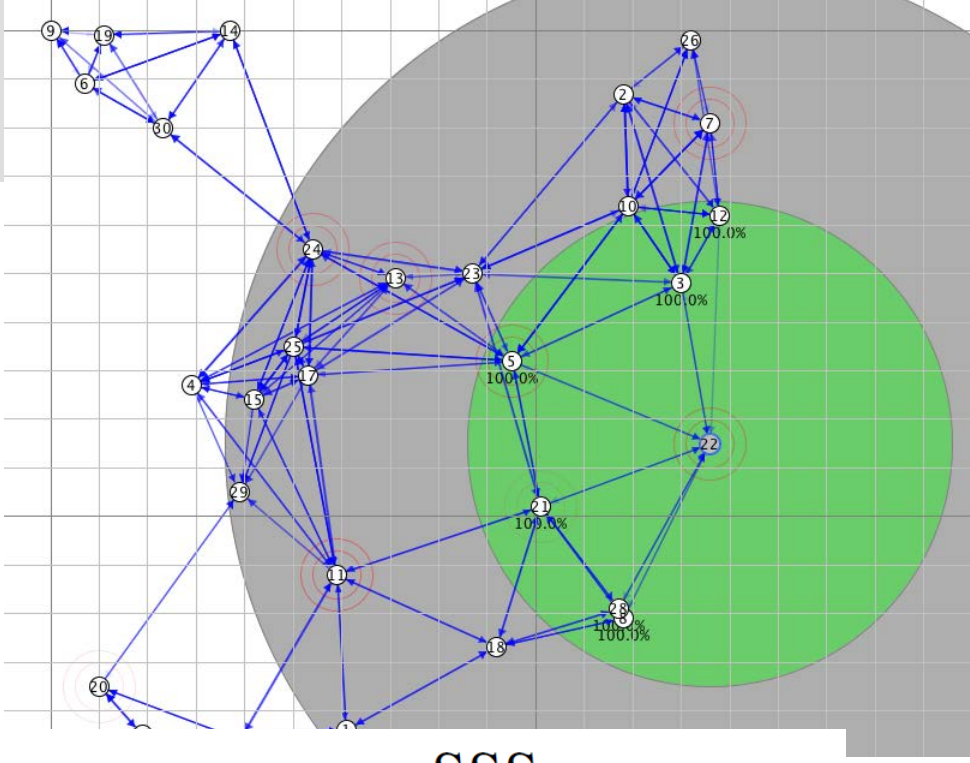
Attack Models – Internal Attacker

- Stolen external storage
  - Key on mote (AES) or overwritten (OTP) ✓
- Storage exhaustion
  - Files for each *$mote_id* limited ✓
- Pollution/data dropping
  - (k; n)-threshold for SSS → < n-k motes ✓
- Stolen/destroyed mote
  - redundancy layer prevents data loss ✓
- Compromised AES key
  - Decryption of future communication
  - Past data can only be recovered when k motes are stolen



Wildlife Monitoring

- 20 of 30 motes have storage
- Constantly, randomly moving
- Different configurations for k



		SSS		
	full redundancy	$k = 2$	$k = 3$	$k = 4$
mean	96.5	88.0	72.8	52.9
σ	3.9	12.2	21.1	25.7

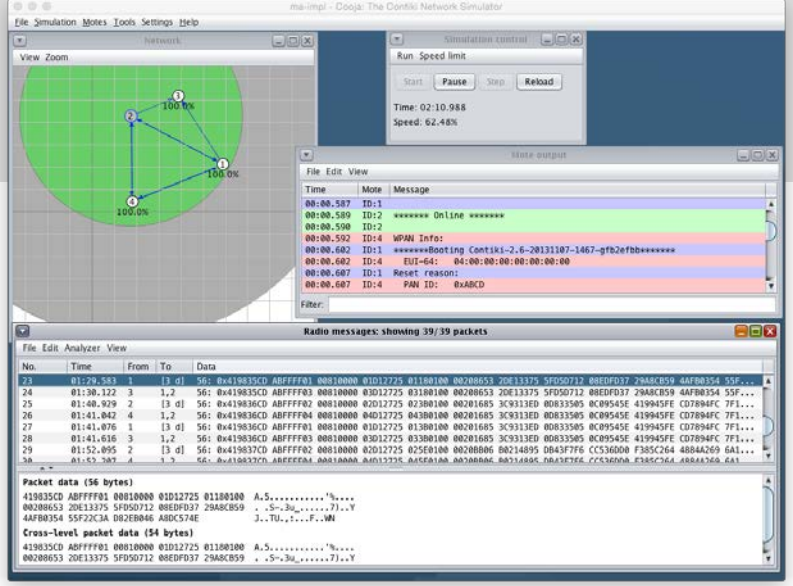
Percentage of recoverable data

- after the loss of one node at different levels of redundancy

Simulations in Cooja

Personal Health Monitoring

- 8 motes randomly generate data
- 4 motes have storage
- Other 4 just send data
- Constantly in radio range
- k = 3
- AES and CBC-MAC



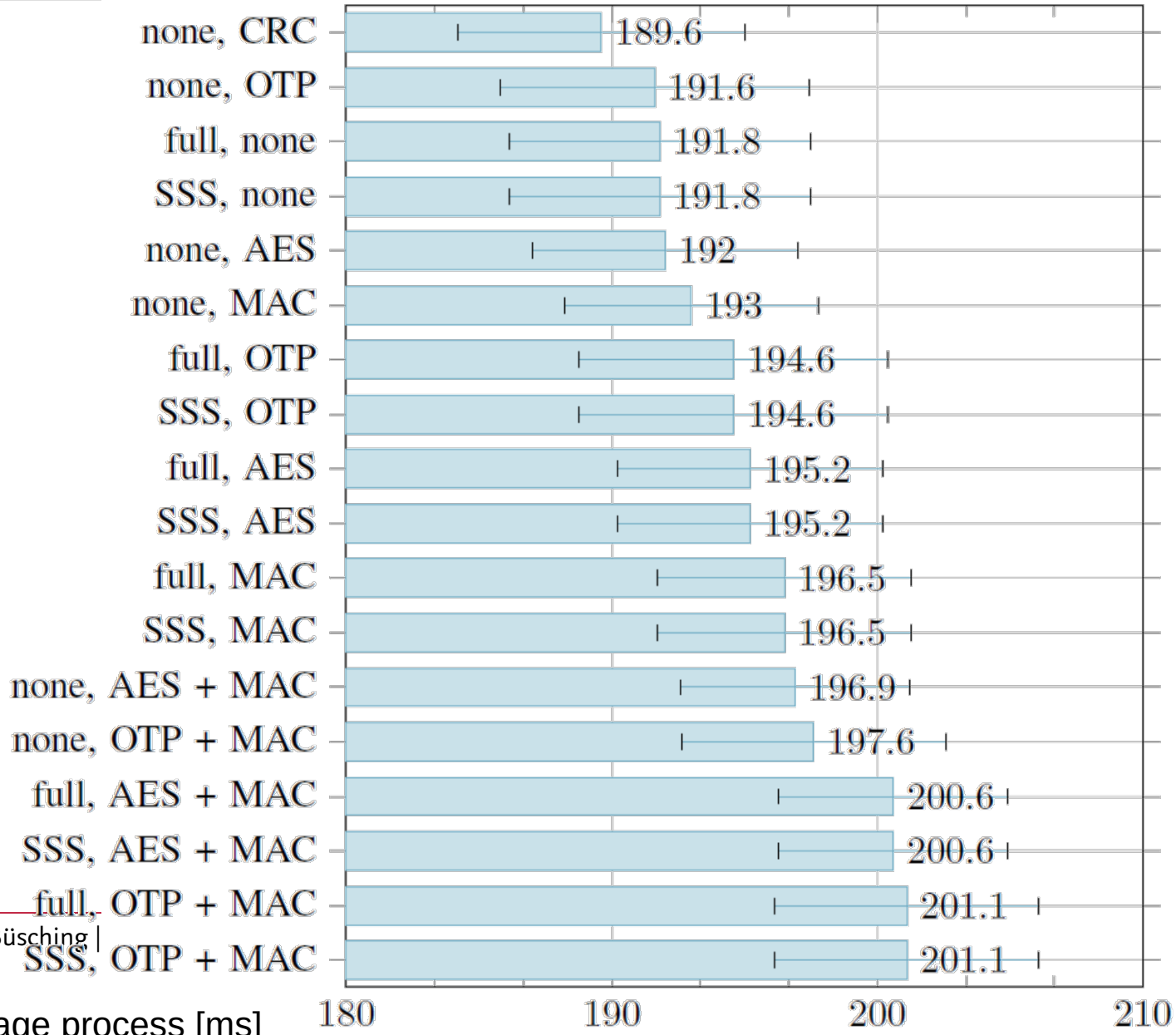
Recoverable data after the loss of one node

mote ID	stored packets	restored packets	[%]
1	118	115	97.5
2	113	112	99.1
3	125	125	100.0
4	119	117	98.3
5	105	105	100.0
6	109	109	100.0
5	118	118	100.0
8	135	134	99.3

with storage

33% storage overhead

Througput (Redundancy, Chipher Suite)



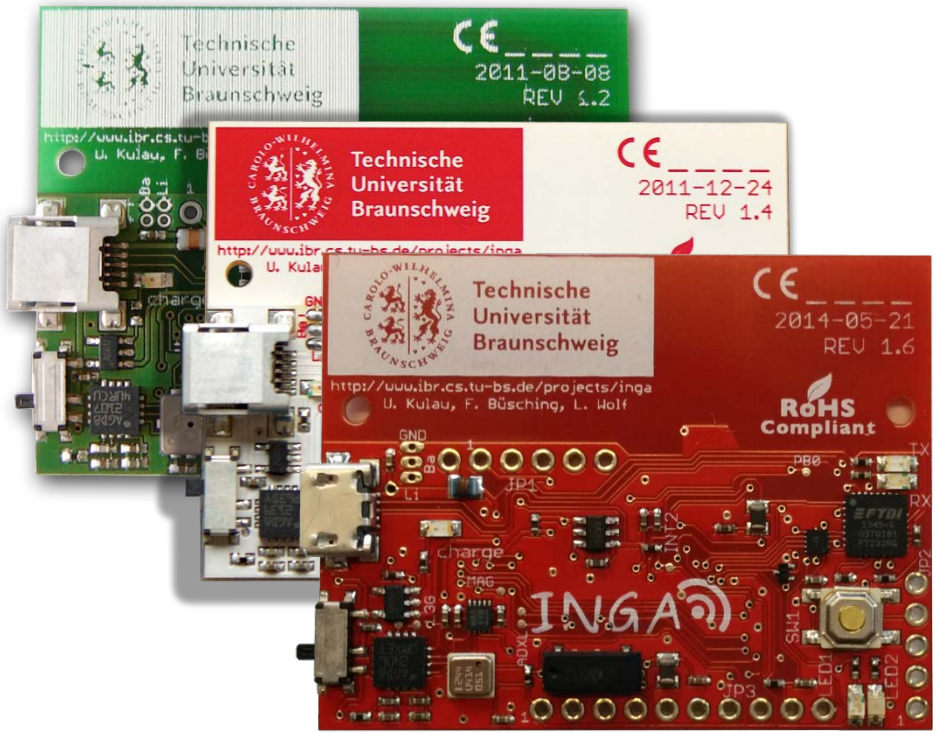
Measurements on physical Motes: INGA

ATmega architecture

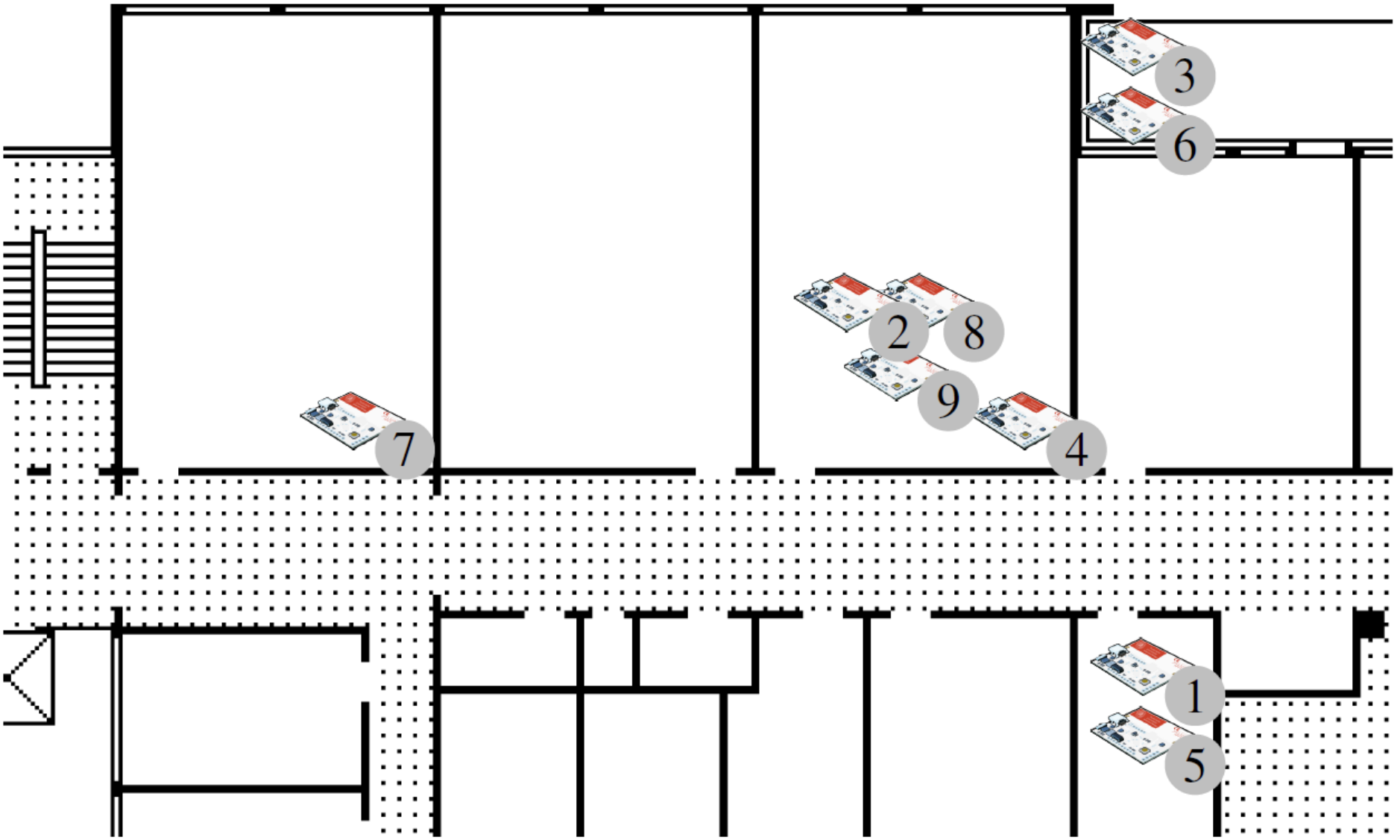
- ATmega 1284p microcontroller
  - 8 bit RISC architecture , 128 kB Flash, 16 kB SRAM, 4 kB EEPROM
- AT86RF233 Radio Transceiver

Setup

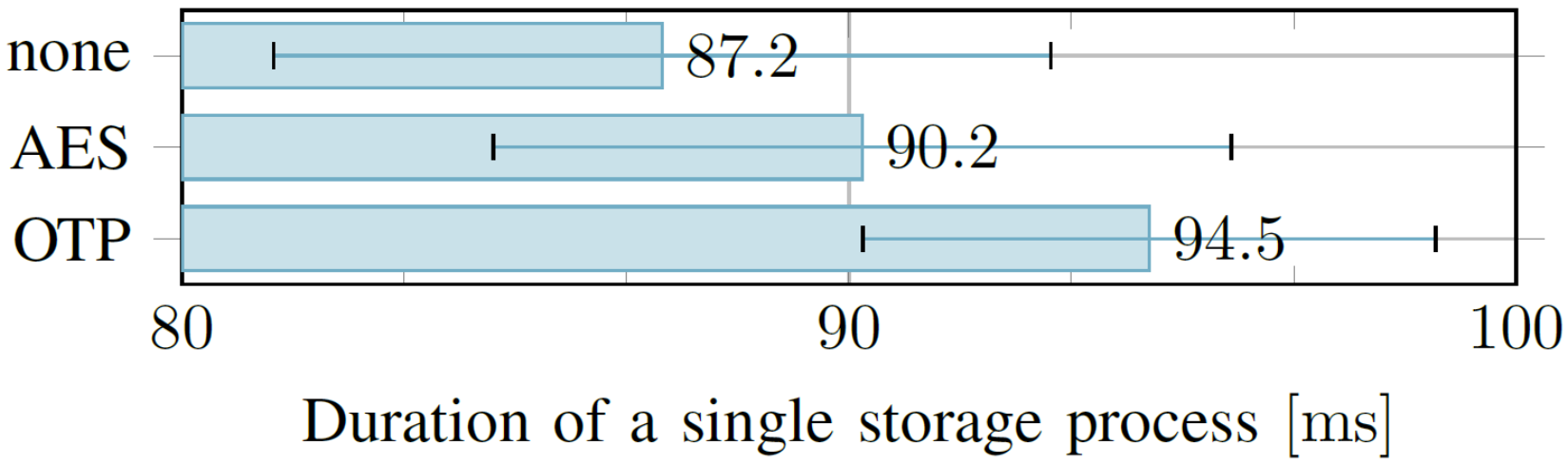
- 9 Motes measure temperature, barometric pressure every 60s



Evaluation in Testbed

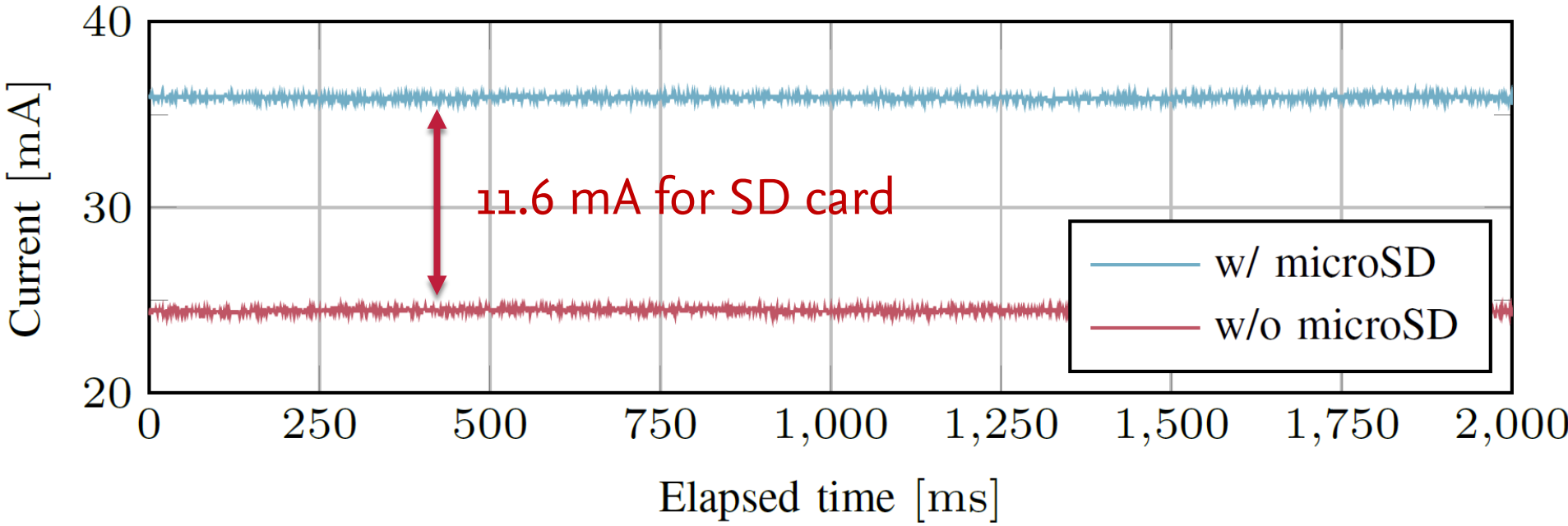


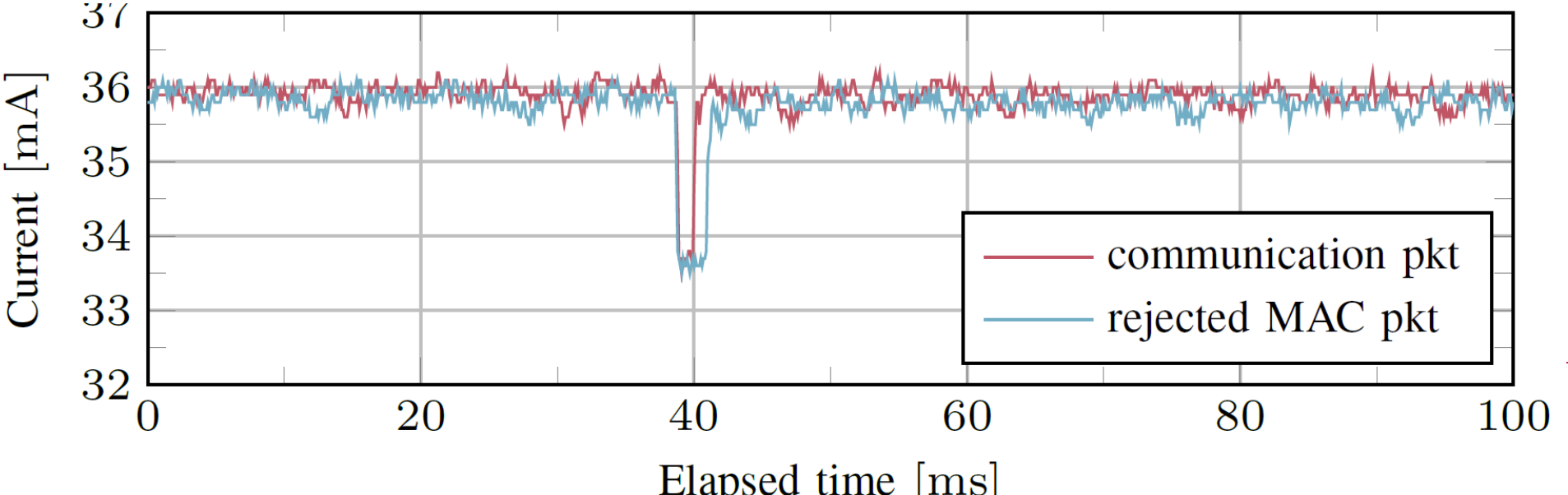
Measured throughput on physical motes



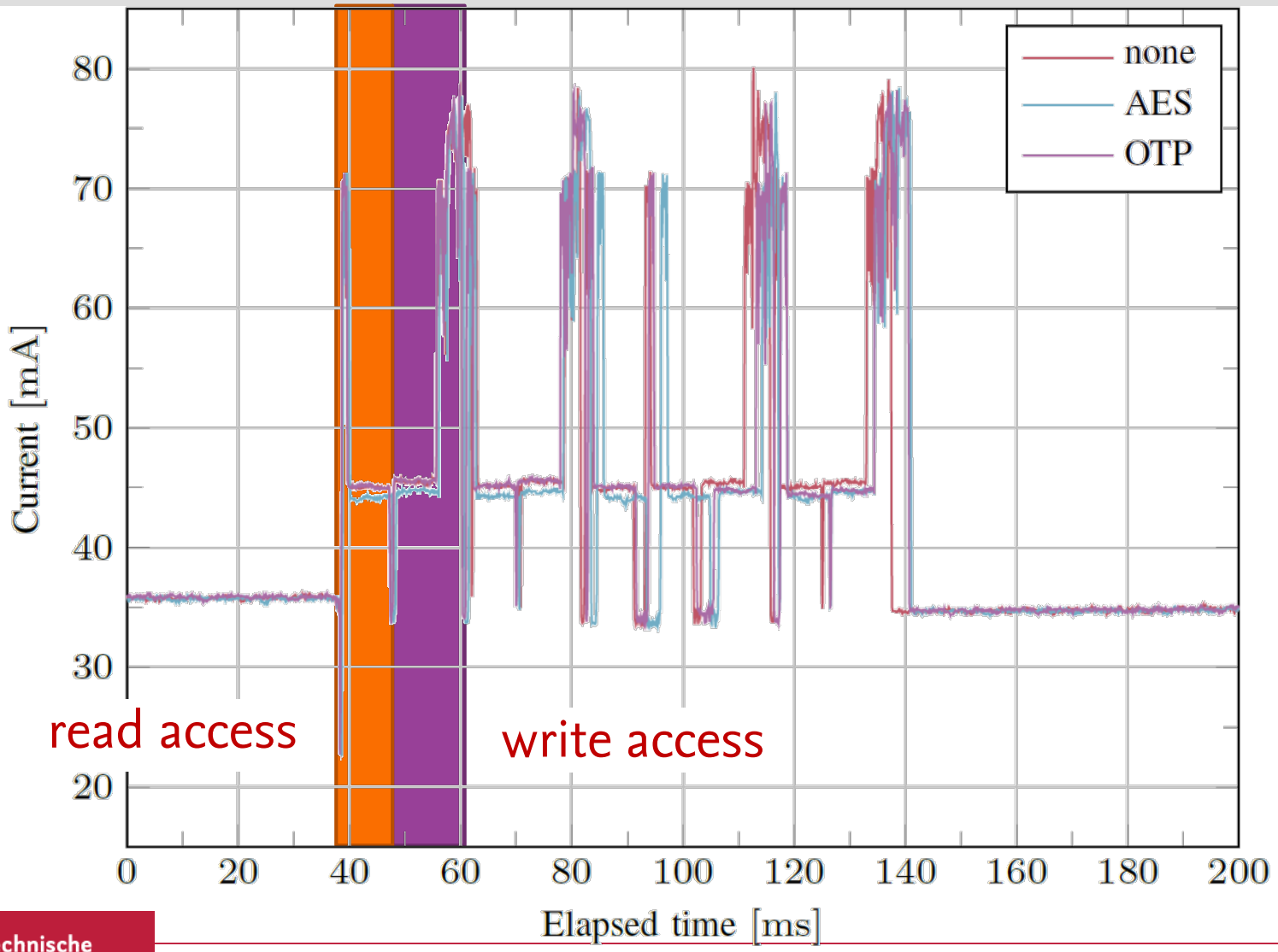
100 ms faster than simulations!

Energy Consumption





Energy Consumption



Summary

We combined
- Shamir's Secret Sharing
- CBC-MAC
- AES / OTP

... to
- Distribute
- Authenticate
- Encrypt

... to ensure configurable
- Redundancy
- Privacy

Cover the whole bandwidth:
- Form full Redundancy to
- Full Confidentiality

RAIM is open soure

check it out at

https://www.ibr.cs.tu-bs.de/projects/raim

Thank you!