

UPPER AND LOWER BOUNDS OF LONG PATHS IN LINE ARRANGEMENTS



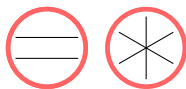
UDO HOFFMANN
TU Berlin

LINDA KLEIST
TU Berlin

TILLMANN MILTZOW
FU Berlin

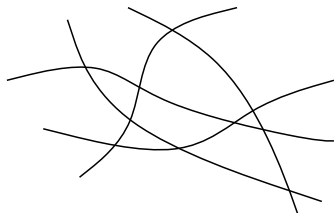
PSEUDOLINE ARRANGEMENTS

SIMPLE

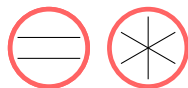


PSEUDOLINE ARRANGEMENTS

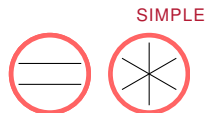
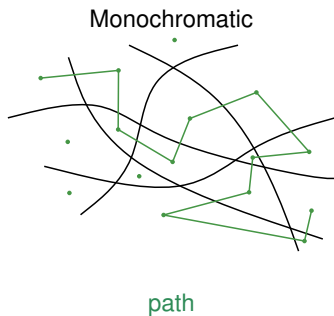
Monochromatic



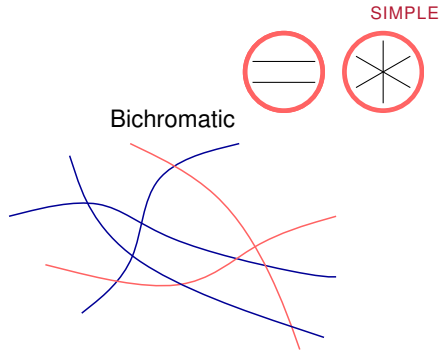
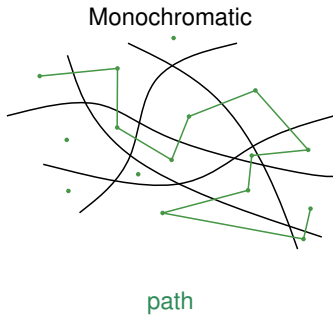
SIMPLE



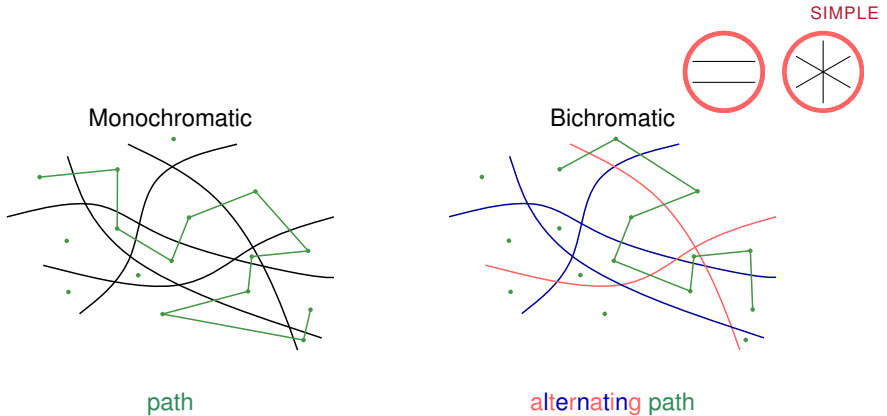
PSEUDOLINE ARRANGEMENTS



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PSEUDOLINE ARRANGEMENTS



MONOCHROMATIC



Aichholzer, Cardinal, Hackl, Hurtado, Korman, Pilz, Silveira,
Uehara, Valtr, Vogtenhuber, Welzl [2014]

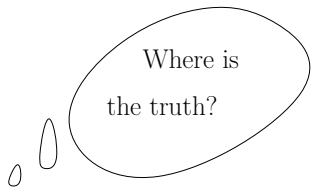
- $\forall \mathcal{A}$ with n lines, longest path $\geq \frac{n^2}{4} + O(n)$
- $\exists \mathcal{A}$ with n lines, longest path $\leq \frac{n^2}{3} + O(n)$

MONOCHROMATIC



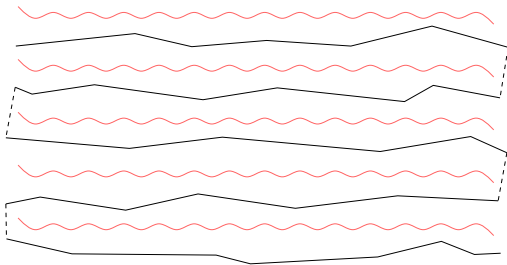
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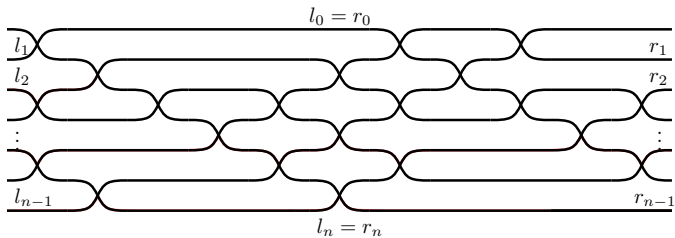
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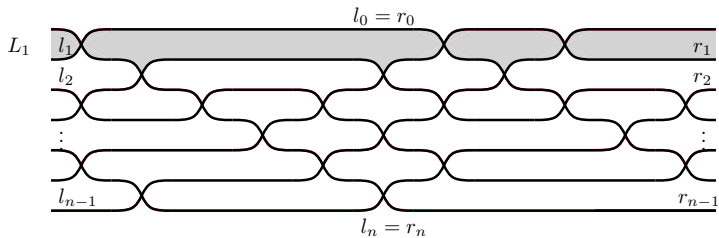


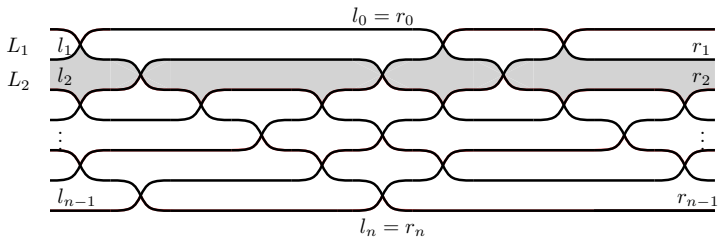
Theorem

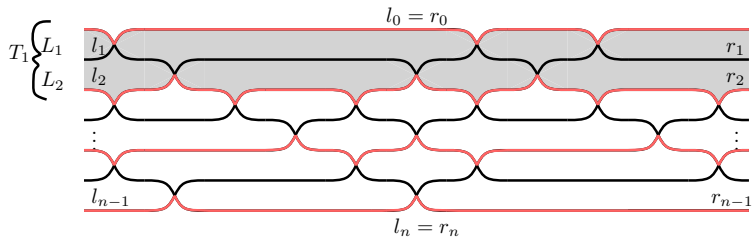
In every simple arrangement \mathcal{A} of n pseudolines, there exists a path of length $\frac{1}{3}n^2 - O(n)$.

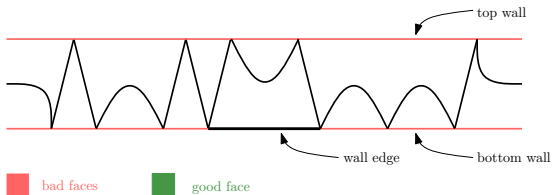
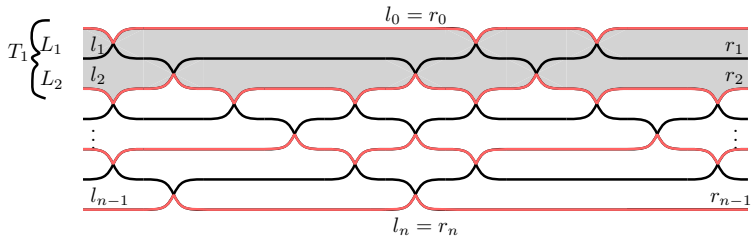


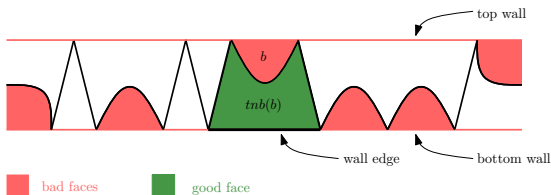
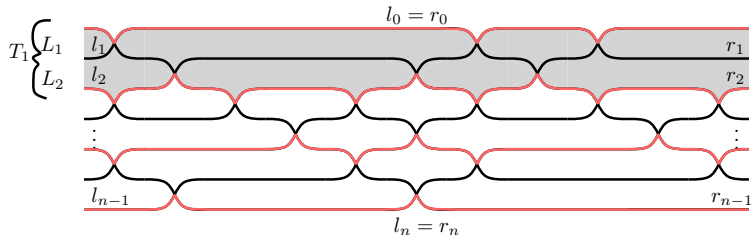


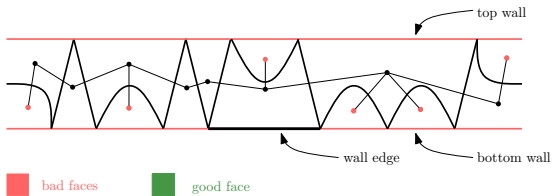
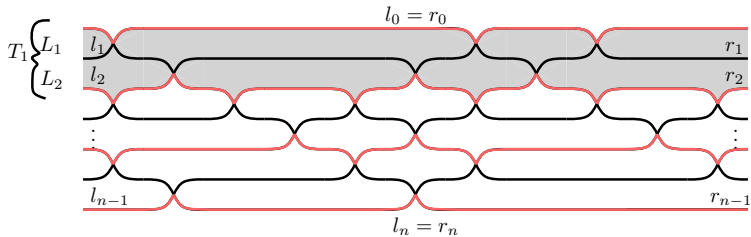


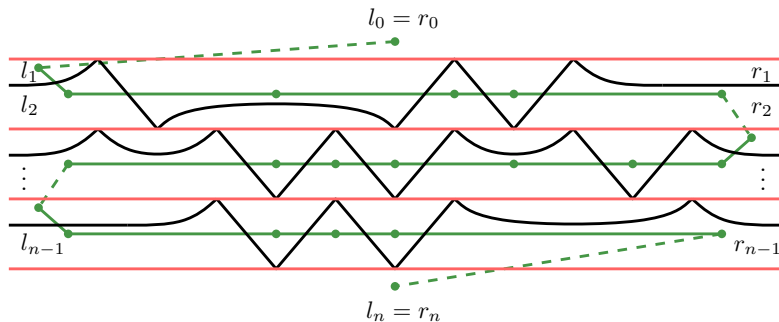












$$P_i := \text{path in tunnel } T_i \begin{cases} \text{from } l_{2i-1} \text{ to } r_{2i} & i \text{ odd,} \\ \text{from } l_{2i} \text{ to } r_{2i-1} & i \text{ even.} \end{cases}$$

$$\sum_{i=1}^{\lfloor \frac{n}{2} \rfloor} |P_i| \geq 1/4 n^2 - n$$

$$F = U + T + N$$

$$ch(f) = \begin{cases} 2 & f \in N, \\ 0 & \text{else.} \end{cases}$$



Find path P and charge function ch such that

- $ch(f) = 0$ for all $f \in N$.
- $ch(f) \leq 1$ for all $f \in T$.
- $ch(f) \leq 2$ for all $f \in U$.
- $\sum_f ch(f) = 2|N|$

	ratio	
2	:	1
traversed	:	not traversed

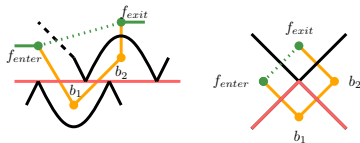
STEP 1: CHARGE-FREENESS FOR N

$p = b_1, b_2, \dots$ maximal path of not traversed faces with $|p| \geq 2$

Step 1a) If $\deg_{\mathcal{A}}(b_1) \geq 3$

reroute,
delete charge,
repeat

$\dots, f_{\text{enter}}, f_{\text{exit}}, \dots \rightsquigarrow \dots, f_{\text{enter}}, b_1, b_2, f_{\text{exit}}, \dots$



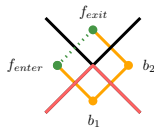
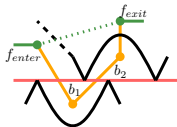
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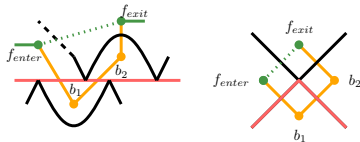
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reroute,
delete charge,
repeat



Step 1b) If $\deg(b_1) = 2$ and $|p| \geq 5$.

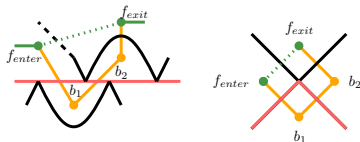
STEP 1: CHARGE-FREENESS FOR N

$p = b_1, b_2, \dots$ maximal path of not traversed faces with $|p| \geq 2$

Step 1a) If $\deg_{\mathcal{A}}(b_1) \geq 3$

reroute,
delete charge,
repeat

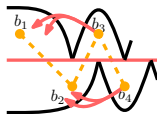
$\dots, f_{\text{enter}}, f_{\text{exit}}, \dots \rightsquigarrow \dots, f_{\text{enter}}, b_1, b_2, f_{\text{exit}}, \dots$



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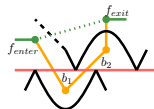
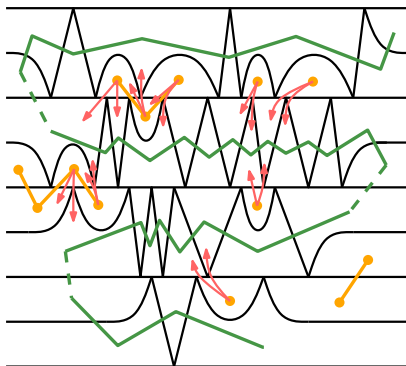
Step 1c) If $\deg(b_1) = 2$ and $|p| \leq 4$.

resend charge



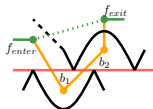
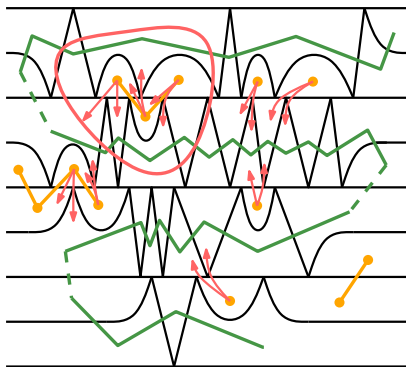
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STEP 1: CHARGE-FREENESS FOR N



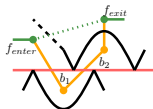
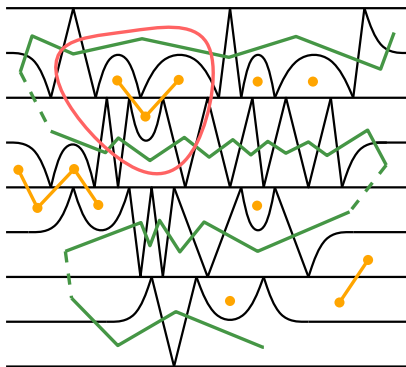
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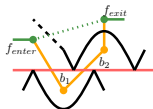
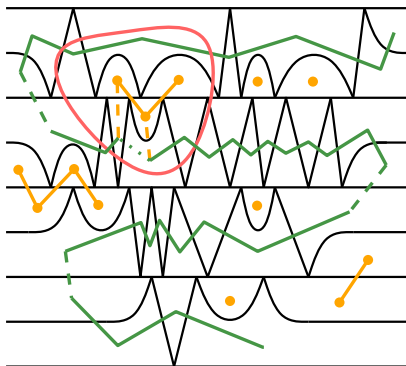
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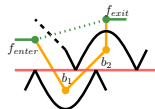
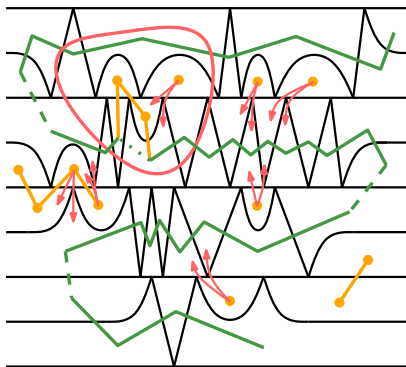
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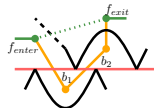
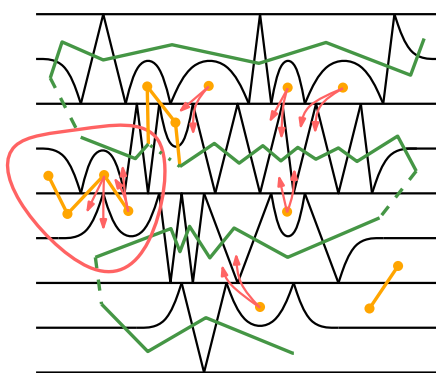
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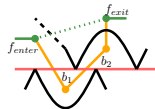
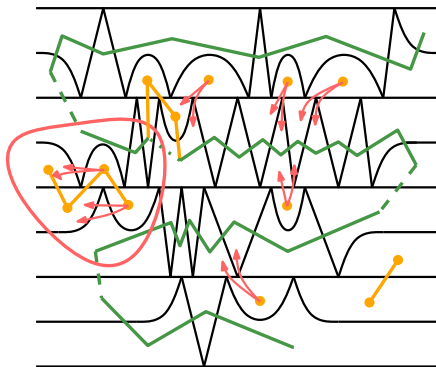
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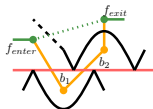
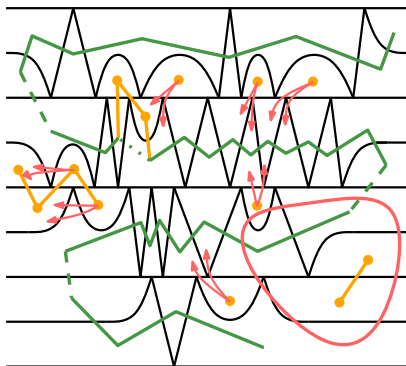
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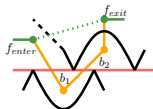
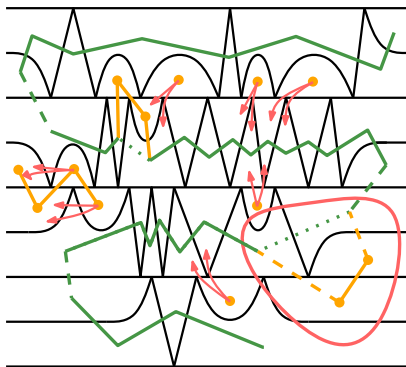
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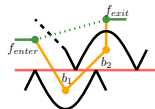
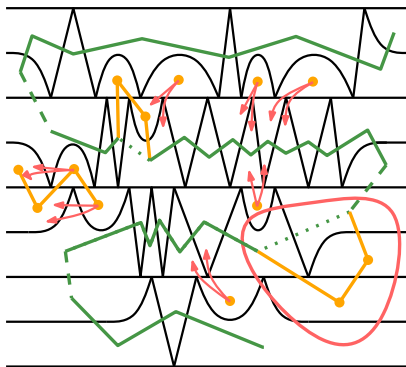
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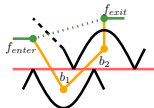
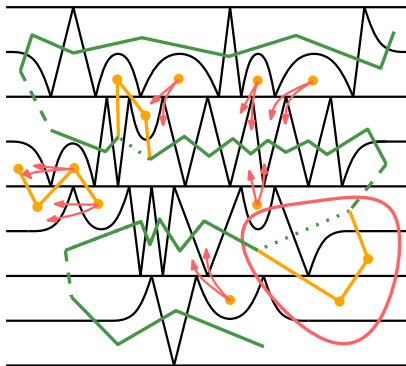


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STEP 1: CHARGE-FREENESS FOR N



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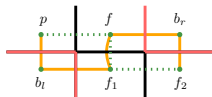
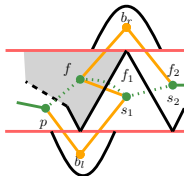
- Not traversed faces are free of charge!
- Only rightmost faces remain sending charge!

STEP 2: ONE UNIT FOR T

f bounded traversed face with charge ≥ 1

Step 2a) If (f_1, f_2) is an edge in P_i $\dots, p, f, f_1, f_2, \dots \rightsquigarrow \dots, p, b_l, f_1, f, b_r, f_2, \dots$

reroute,
delete charge

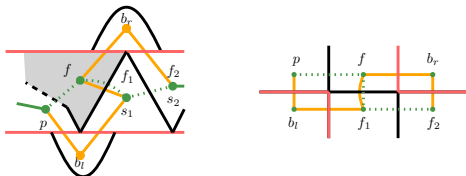


STEP 2: ONE UNIT FOR T

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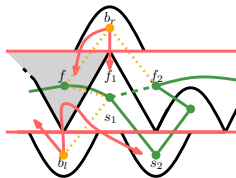
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reroute,
delete charge



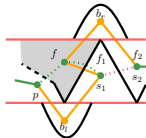
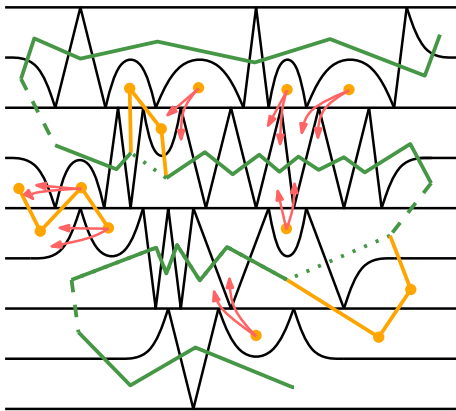
Step 2b) If (f_1, f_2) was replaced in Step 1

resend charge



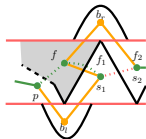
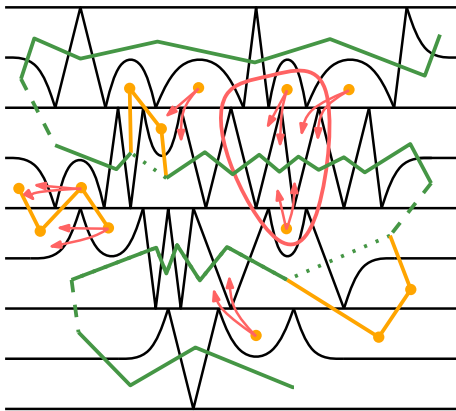
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STEP 2: ONE UNIT FOR T



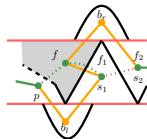
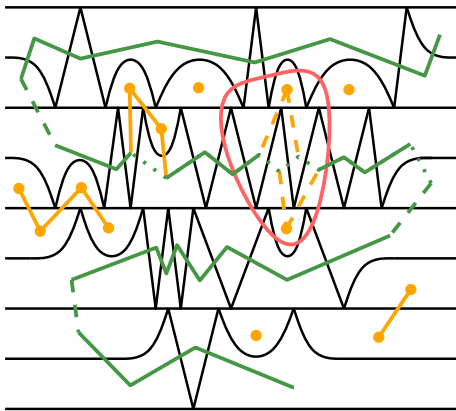
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STEP 2: ONE UNIT FOR T



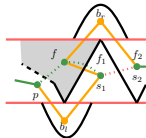
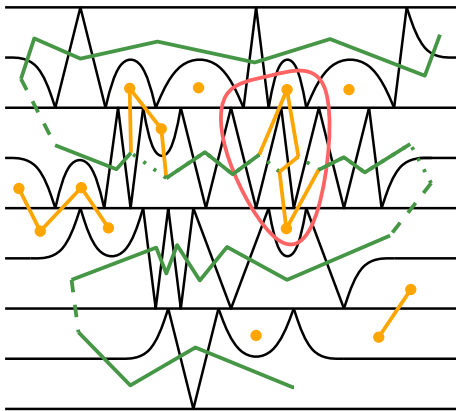
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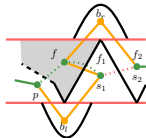
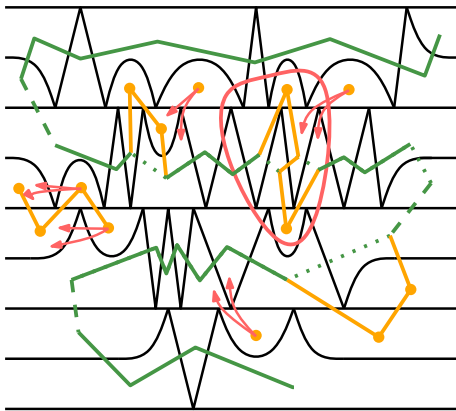
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STEP 2: ONE UNIT FOR T



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We found path P and charge function ch such that

- $ch(f) = 0$ for all $f \in N$.
- $ch(f) \leq 1$ for all $f \in T$.
- $ch(f) \leq 2$ for all $f \in U$.
- $\sum_f ch(f) = 2|N|$

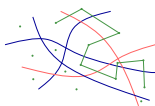


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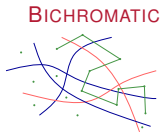
BICHROMATIC



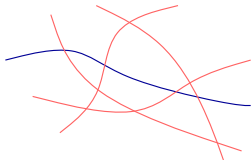
Aichholzer, Cardinal, Hackl, Hurtado, Korman, Pilz, Silveira,
Uehara, Valtr, Vogtenhuber, Welzl [2014]

- \forall bicolored \mathcal{A} , longest alternating paths $\geq n$
- \exists bicolored \mathcal{A} with longest path $\leq 2n + 1$

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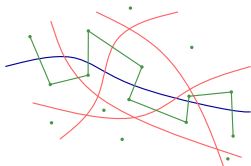
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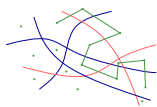


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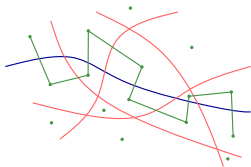


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Uehara, Valtr, Vogtenhuber, Welzl [2014]

BICHROMATIC



- \forall bicolored \mathcal{A} , longest alternating paths $\geq n$
- \exists bicolored \mathcal{A} with longest path $\leq 2n + 1$



What about
arrangements with
 n red and
 n blue lines?

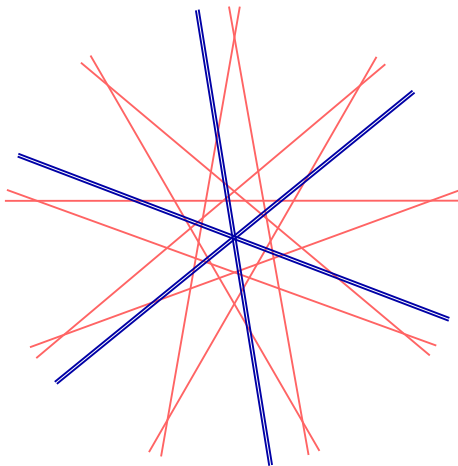
Does every arrangement of n red and n blue lines has an alternating path of quadratic length?

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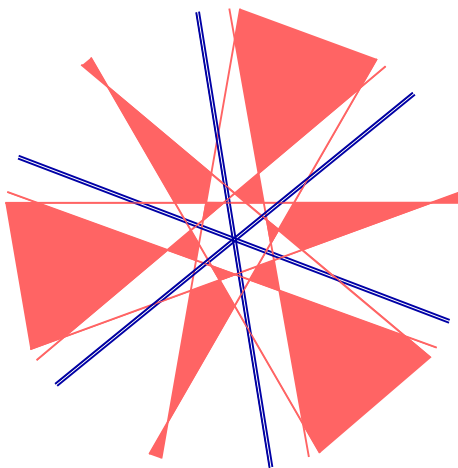
Theorem

There exists a simple arrangement \mathcal{A} of $3k$ red and $2k$ blue lines where any alternating dual path goes through at most $14k$ faces, for every odd k .

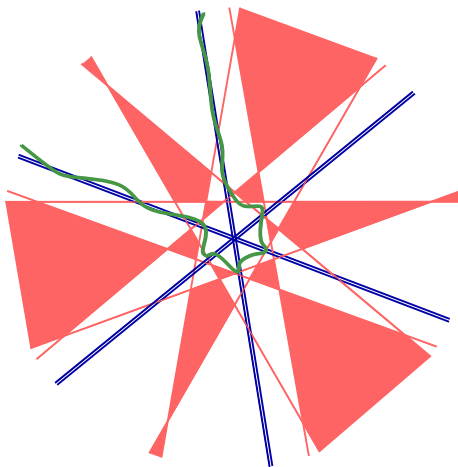
UPPER BOUND EXAMPLE



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Given an arrangement \mathcal{A} of n lines, is there a coloring such that there exist an alternating path of length $\Omega(n^2)$?

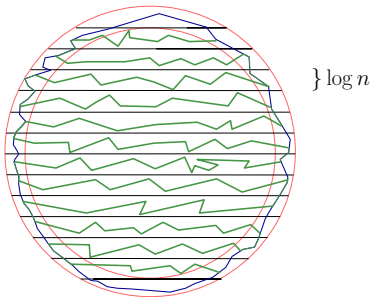
Given an arrangement \mathcal{A} of n lines, is there a coloring such that there exist an alternating path of length $\Omega(n^2)$?

Theorem

In a random bicolored of an arrangement of n pseudolines, there exists an alternating path of length $\Omega(n^2/\log n)$ with high probability.

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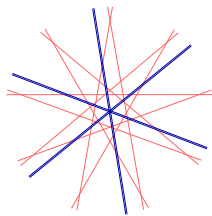
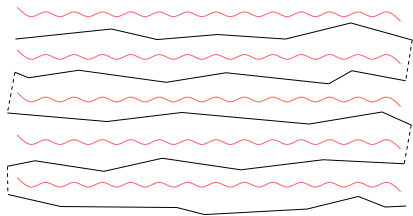
1 Bicolored arrangements:

Gap between lower and upper bound

$$\left[n \leq \min_{\mathcal{A}: |\mathcal{A}|=n} \max_{P \in \mathcal{A}} |P| \leq 2n + 1 \right]$$

2 Balanced Bicolored Arrangements:

In an arrangement of n red and n blue pseudolines, is there an alternating path of quadratic length?



- tight lower bound for monochromatic path $[n^2/3 - O(n)]$
- upper bound for almost balanced bichromatic arrangement $[O(n)$ for $3n$ red and $2n$ blue lines]
- $\Omega(n^2/\log n)$ for a random coloring with high probability

